### Lineage Stash: Fault Tolerance Off the Critical Path

Stephanie Wang, John Liagouris, Robert Nishihara, Philipp Moritz, Ujval Misra, Alexey Tumanov, Ion Stoica



## Low latency is increasingly important in data processing systems

Data processing is used today in **online** systems.

- Stream processing
- Graph processing
- Control systems

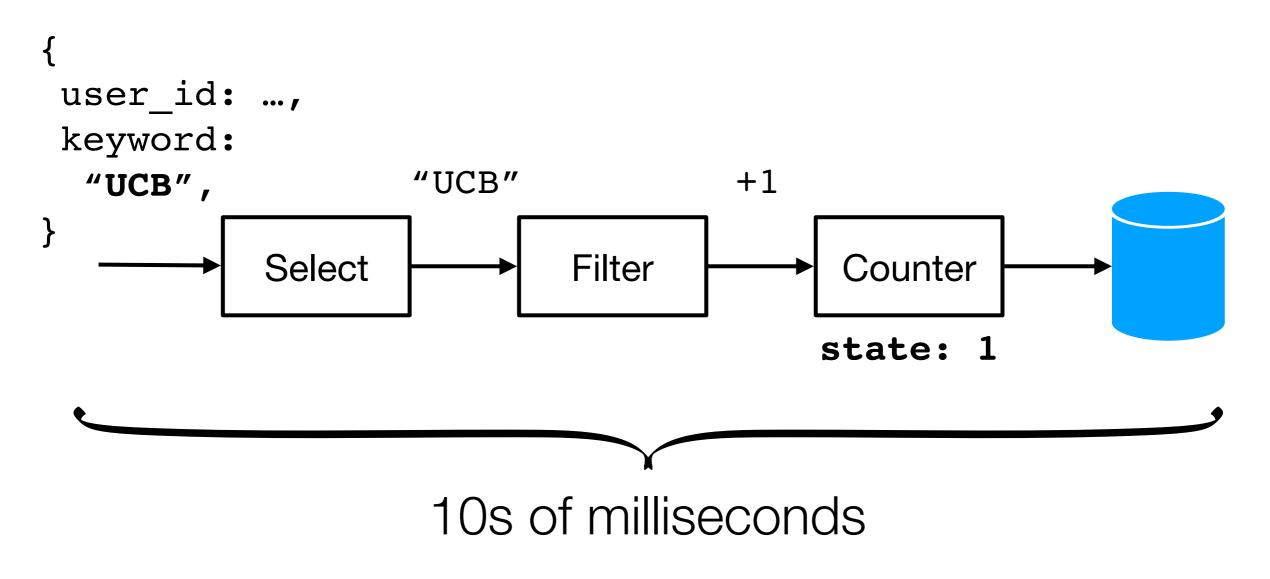
Data processing at **large scale** also requires the ability to **recover** results after a failure.

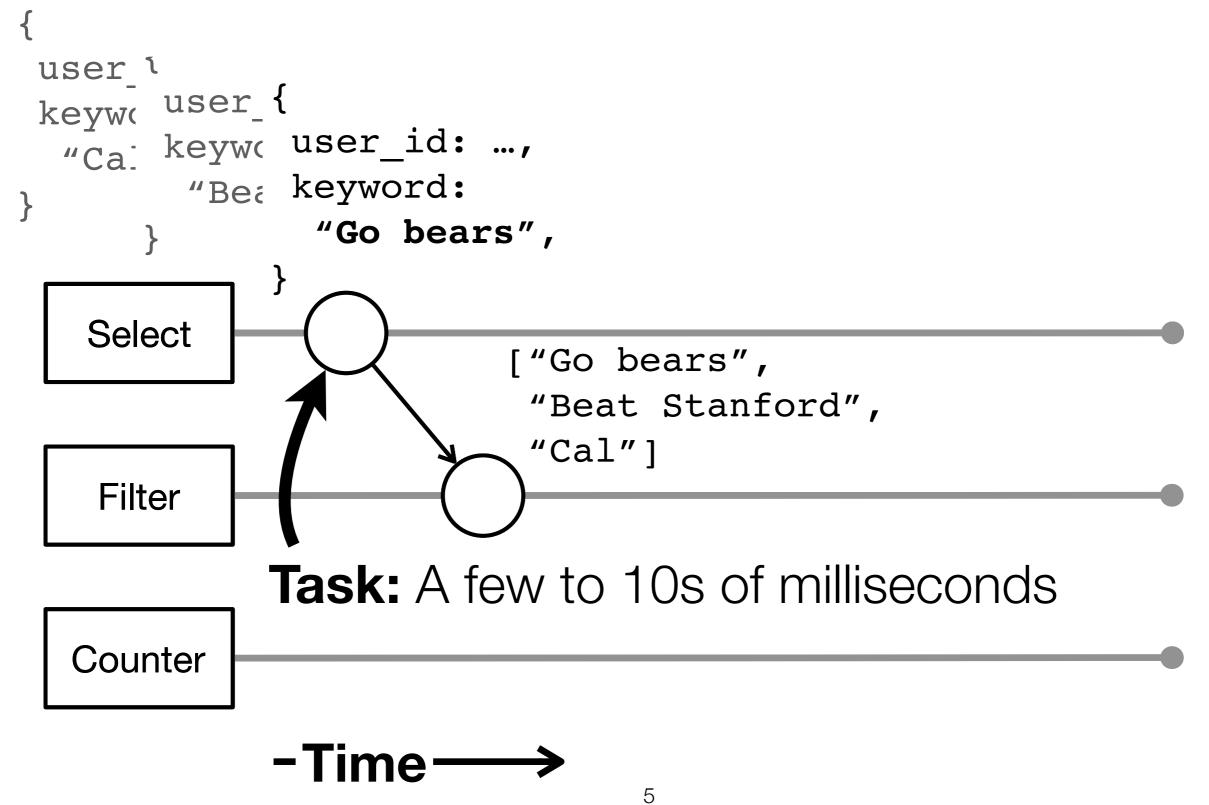
# Tradeoff between low latency and recovery time

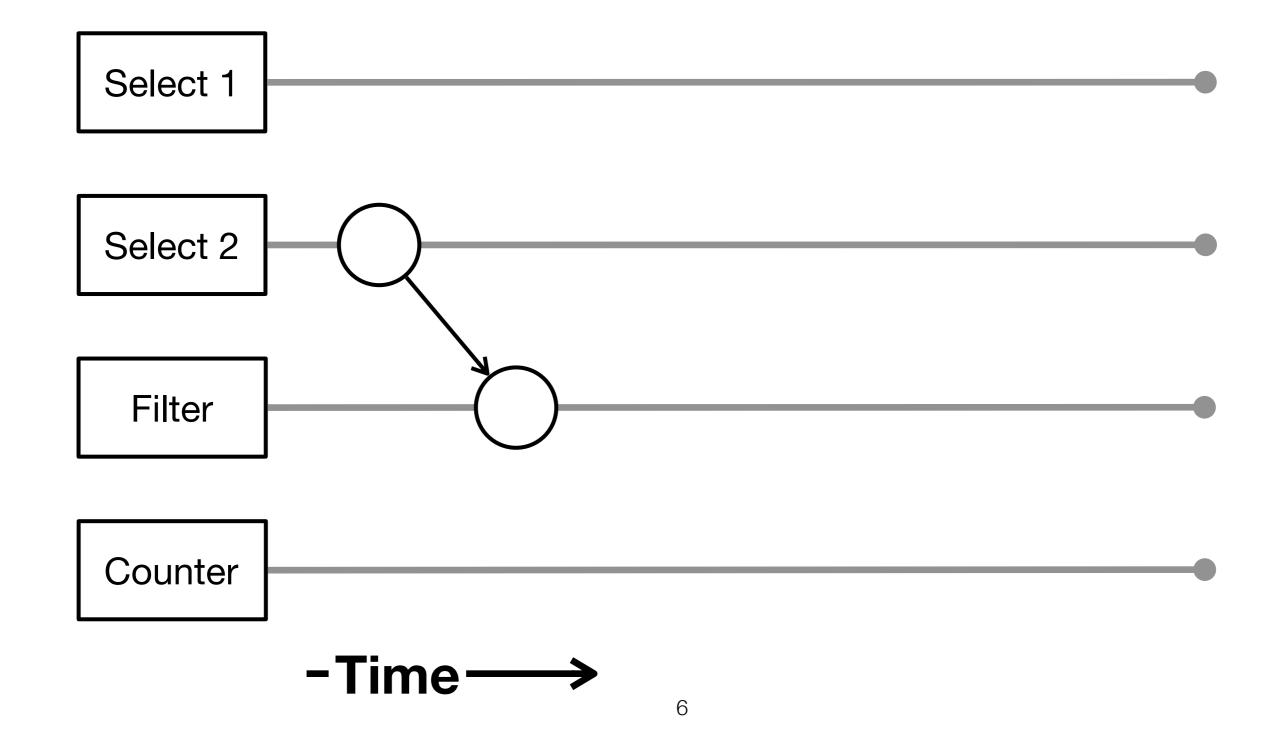


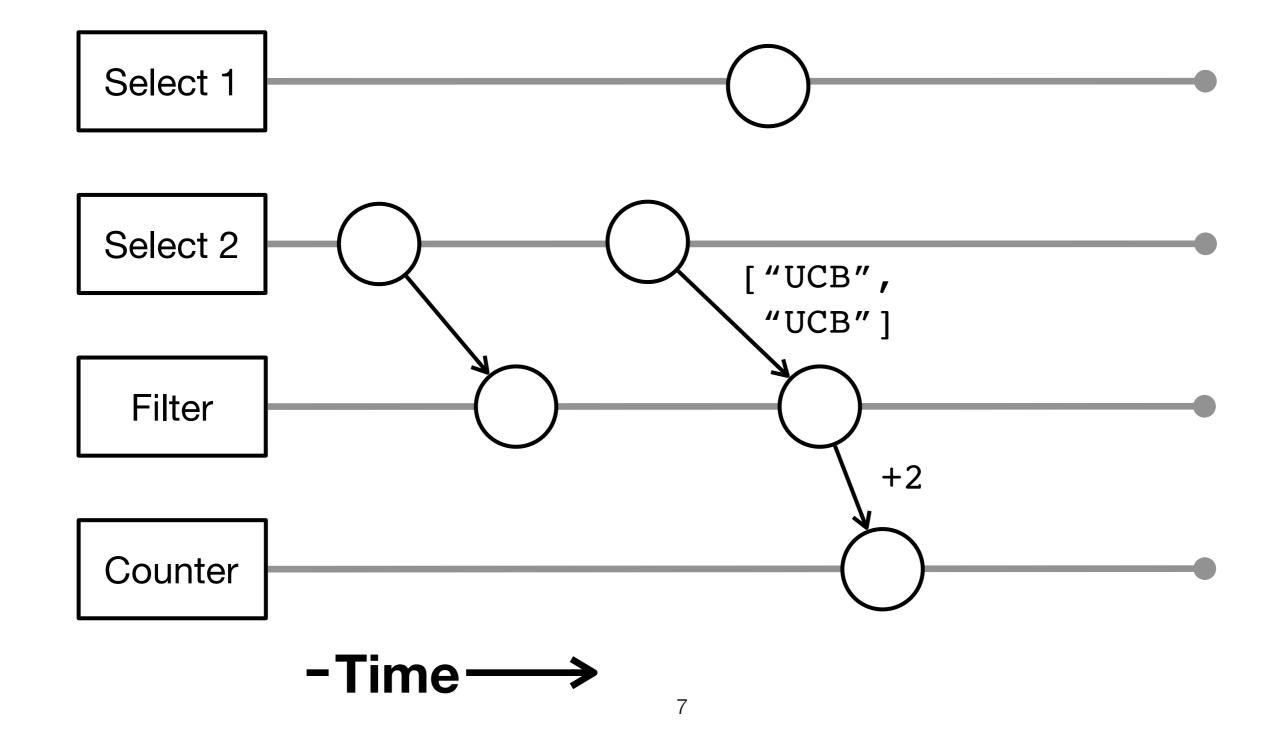
Runtime overhead

Display the number of times "UCB" has been queried.

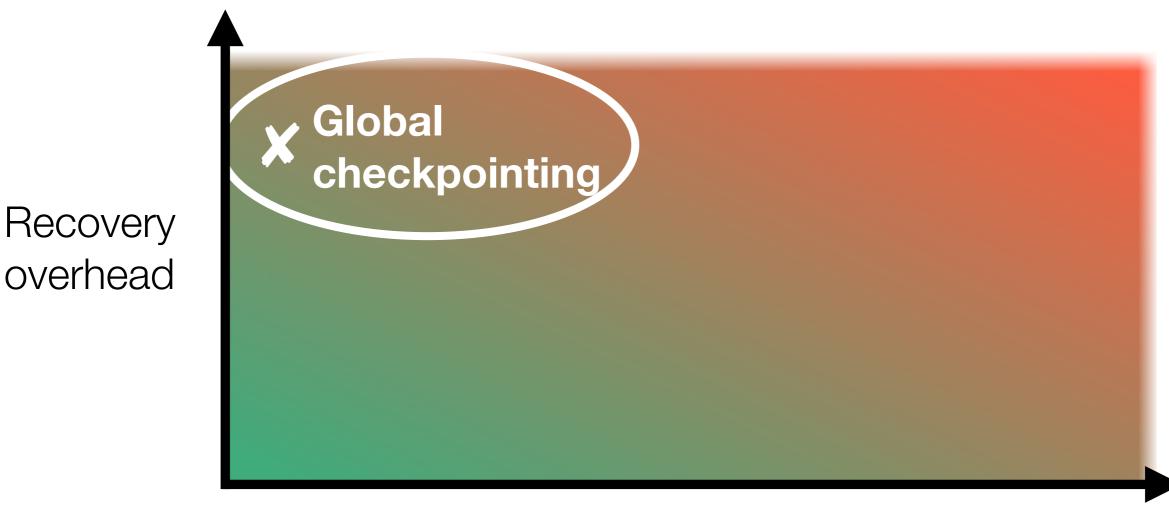




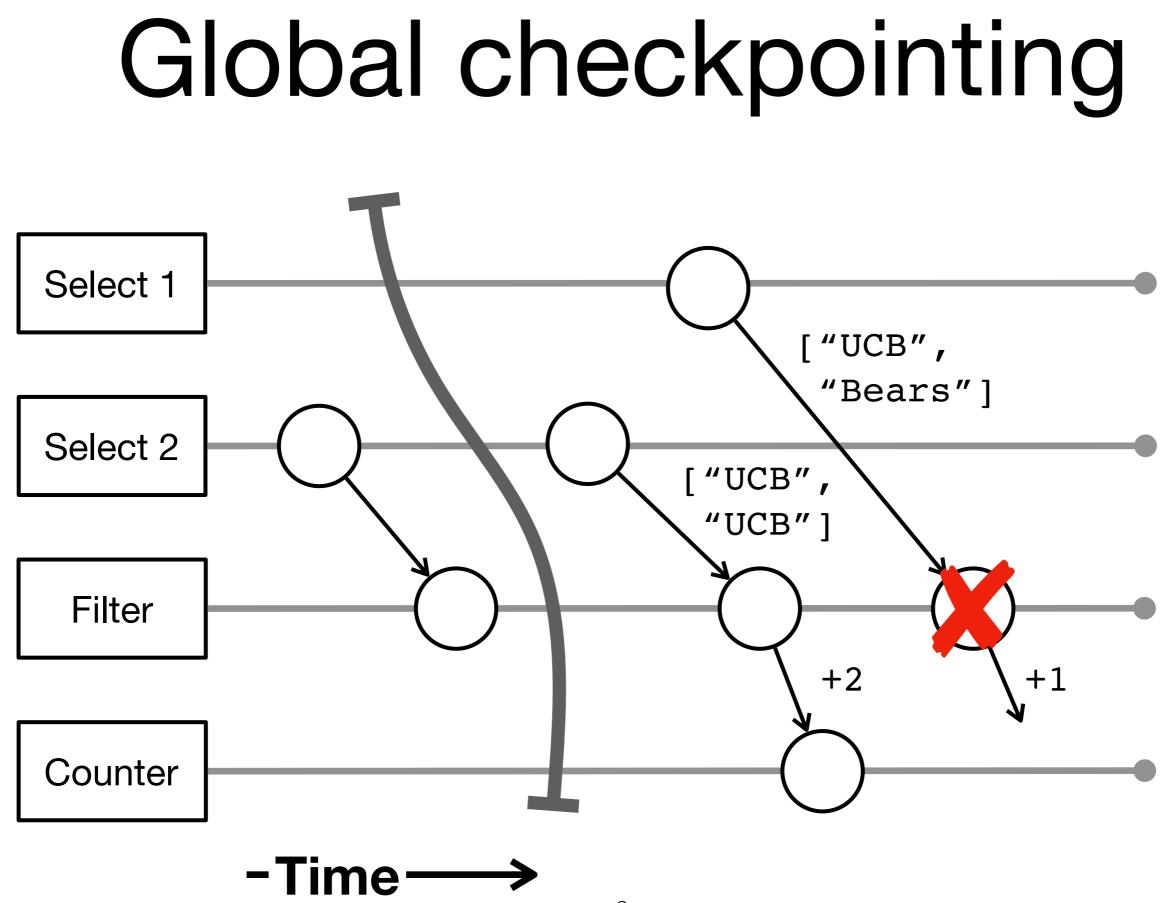


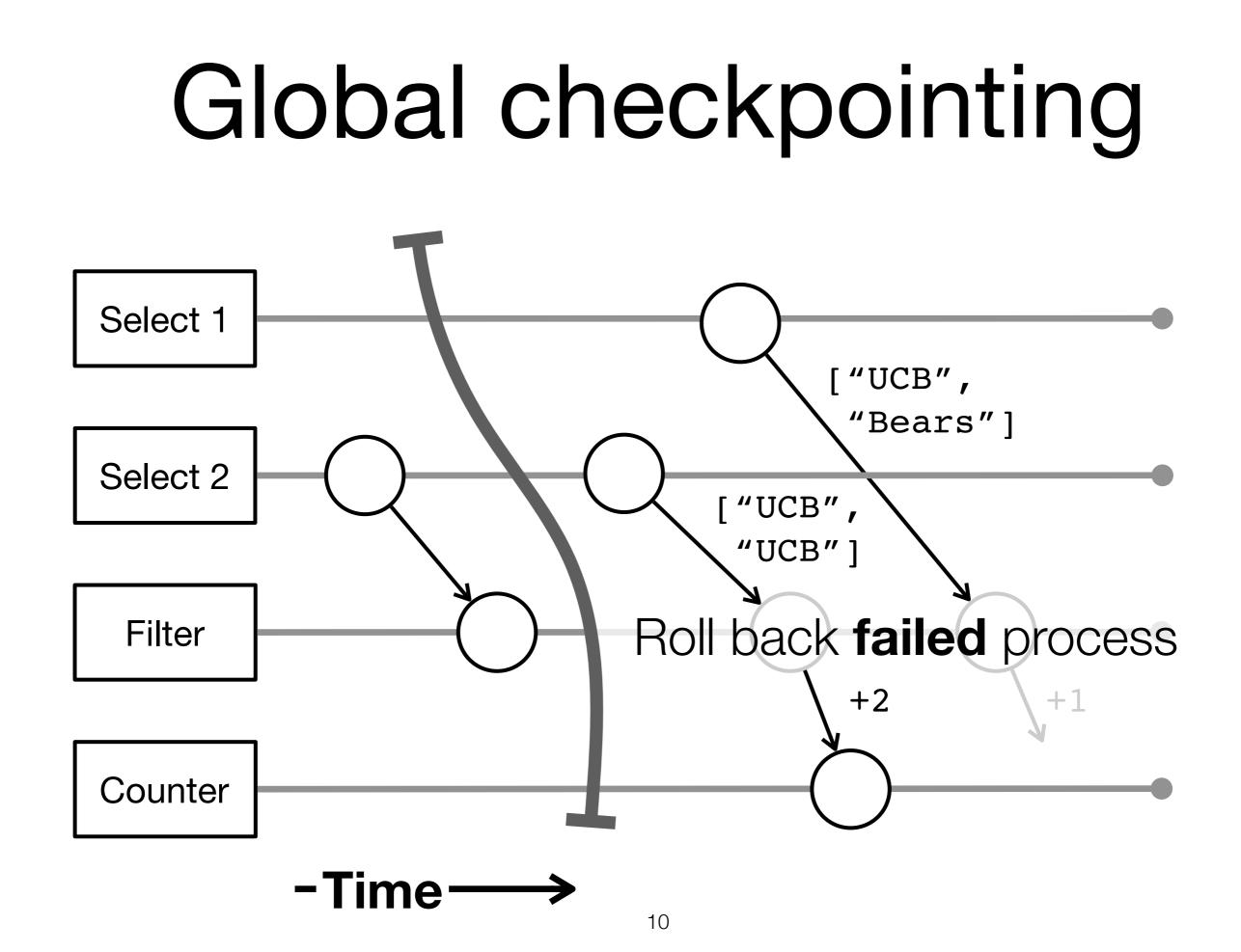


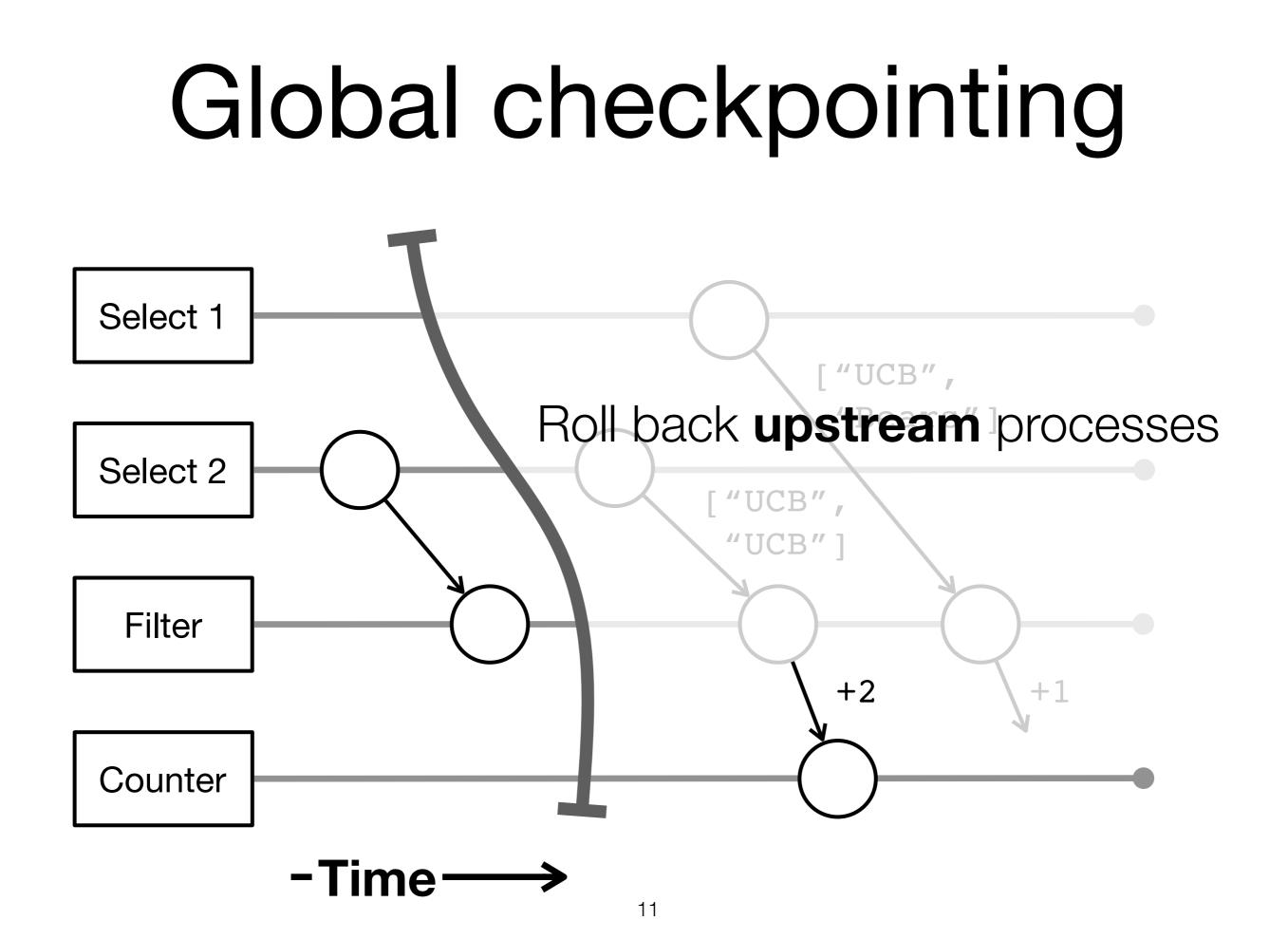
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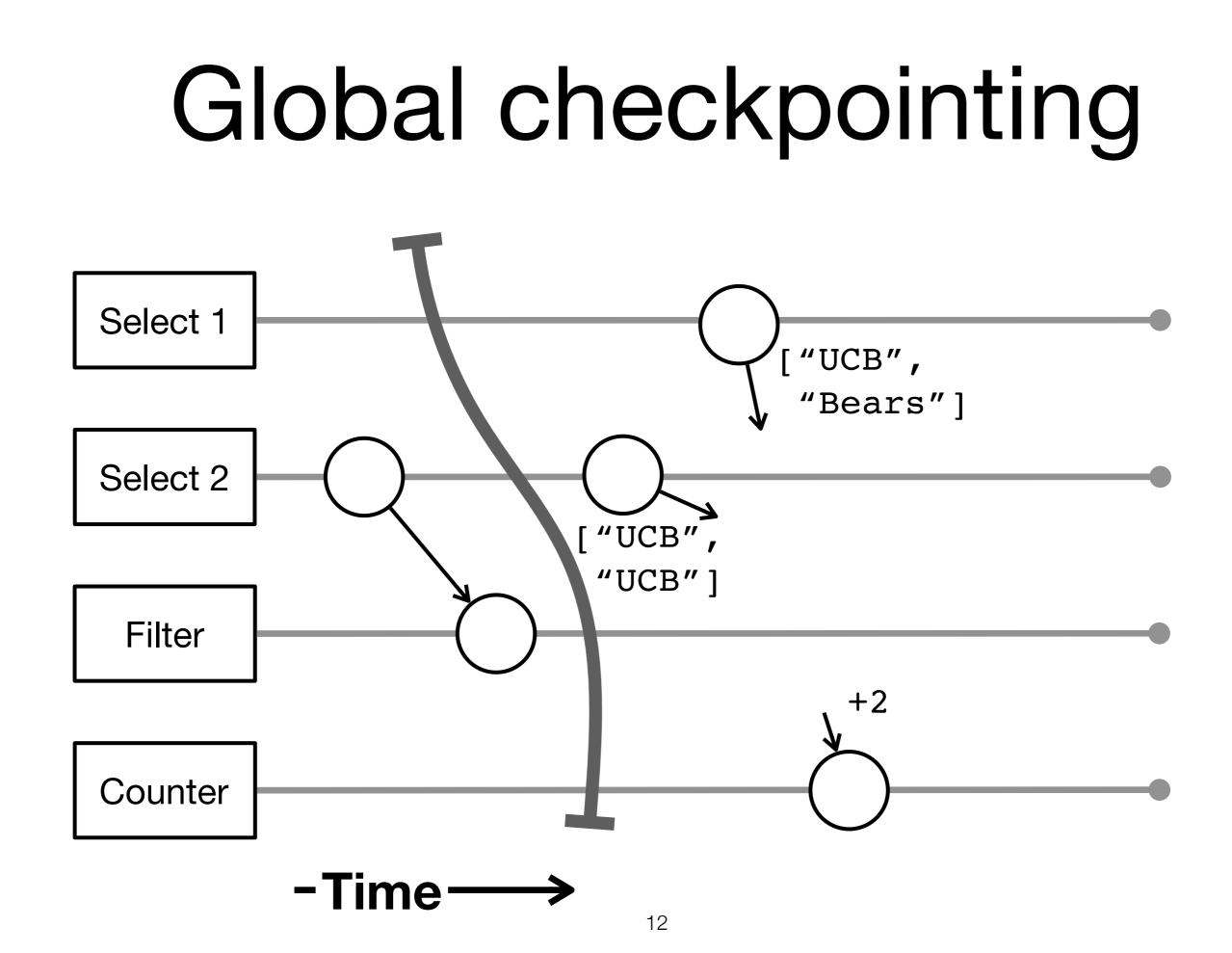


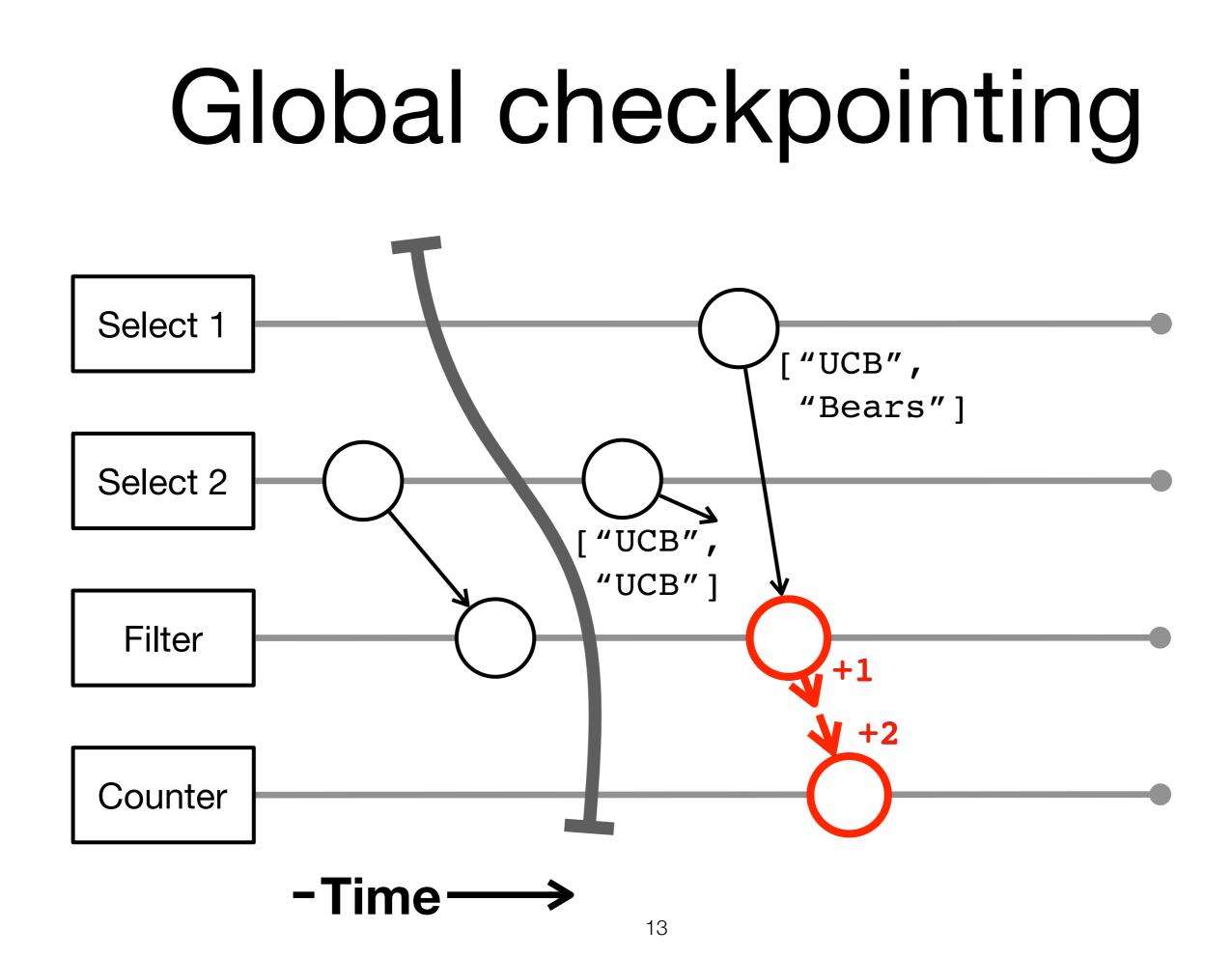
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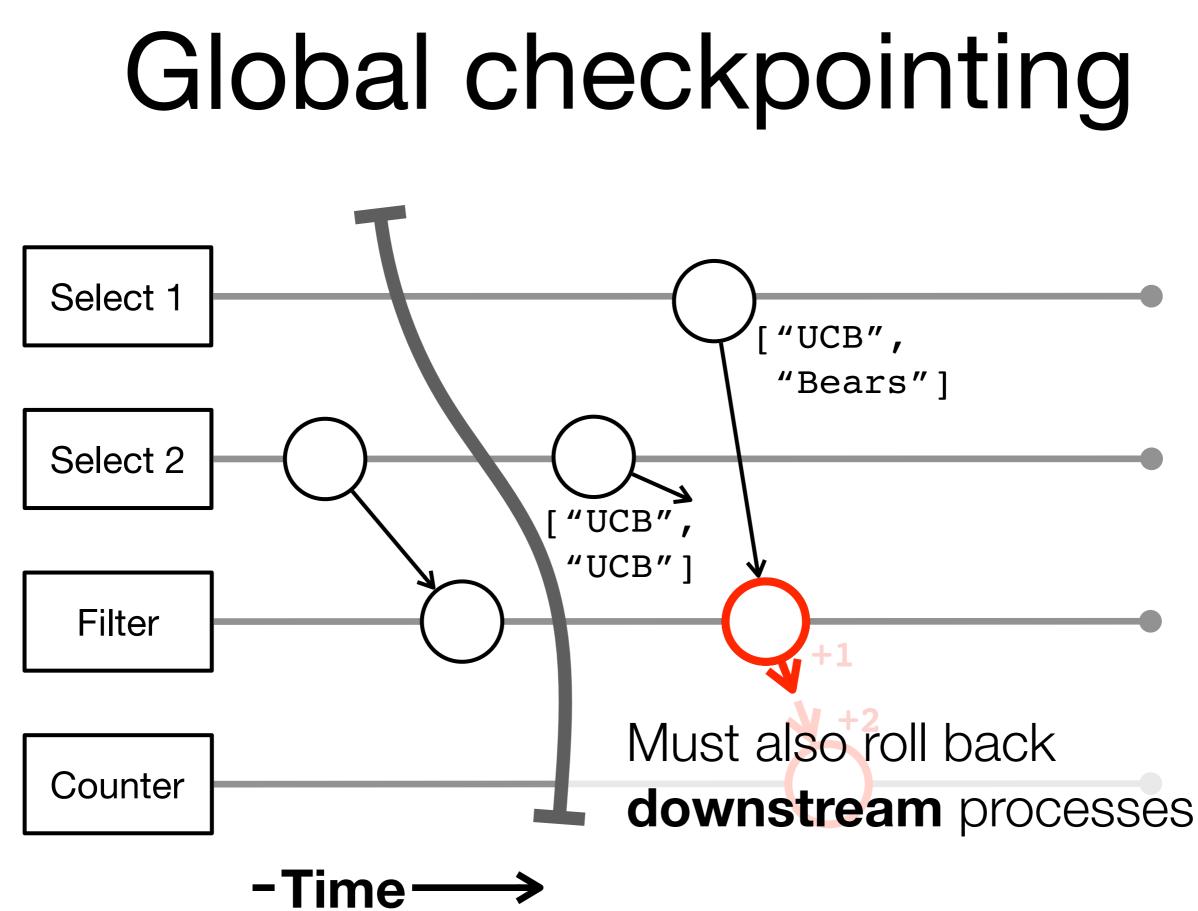


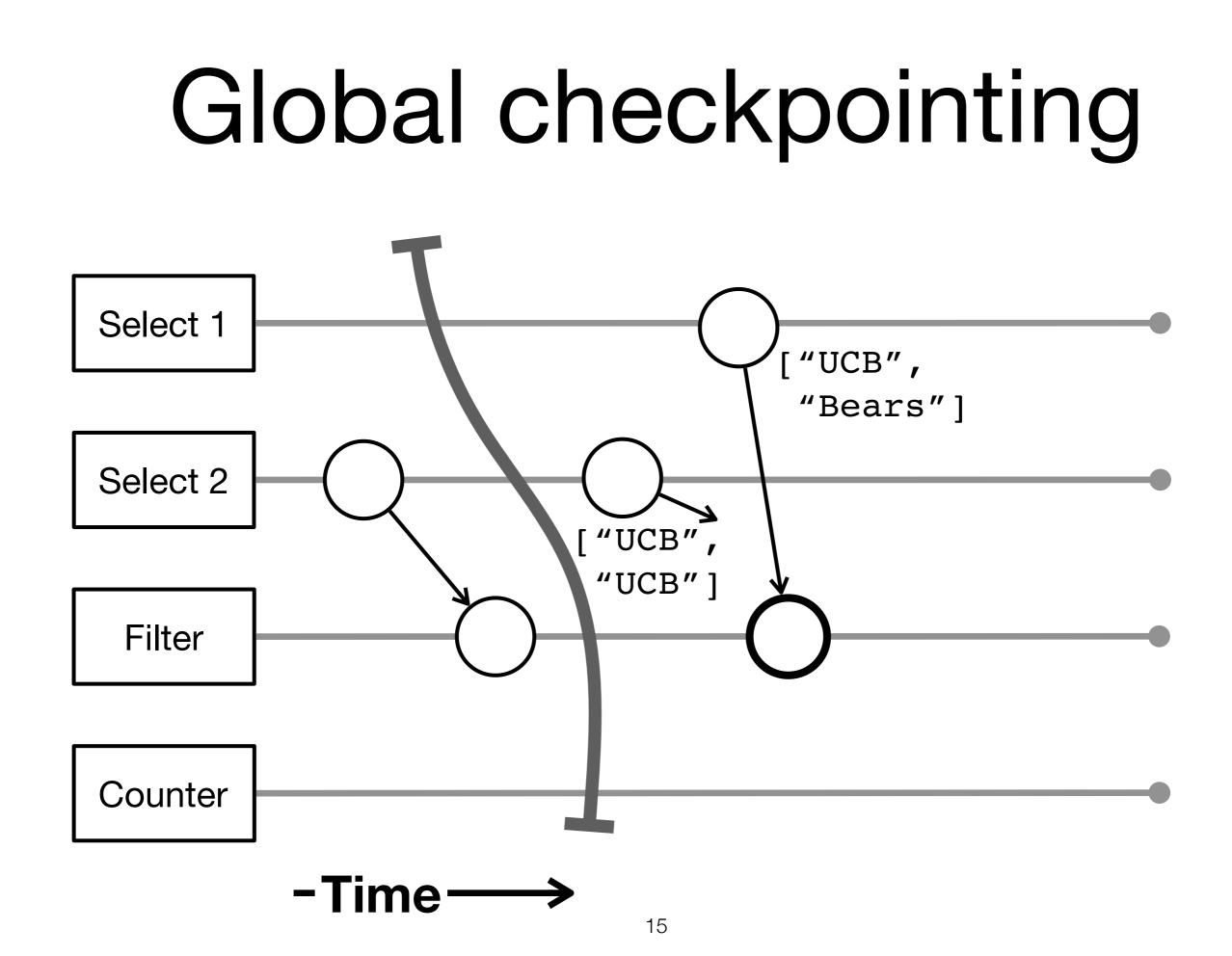


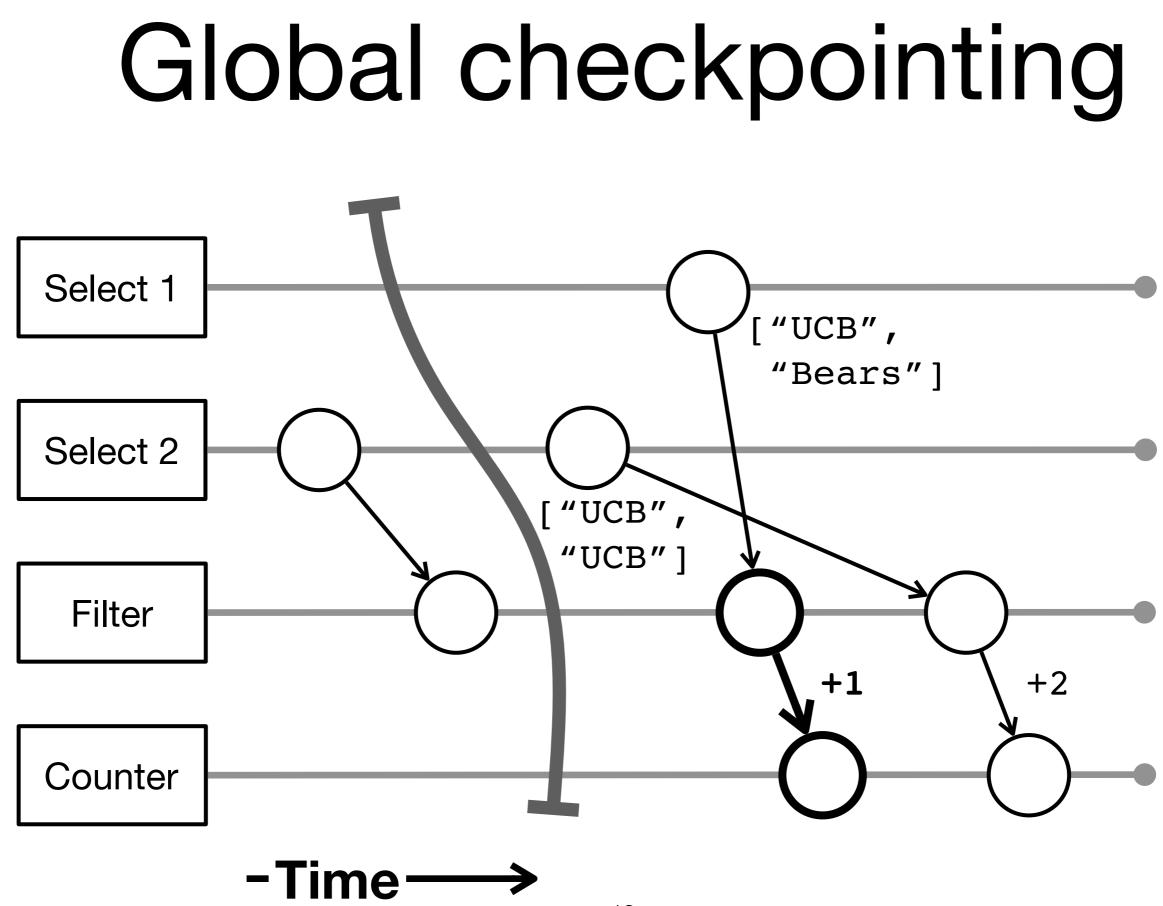


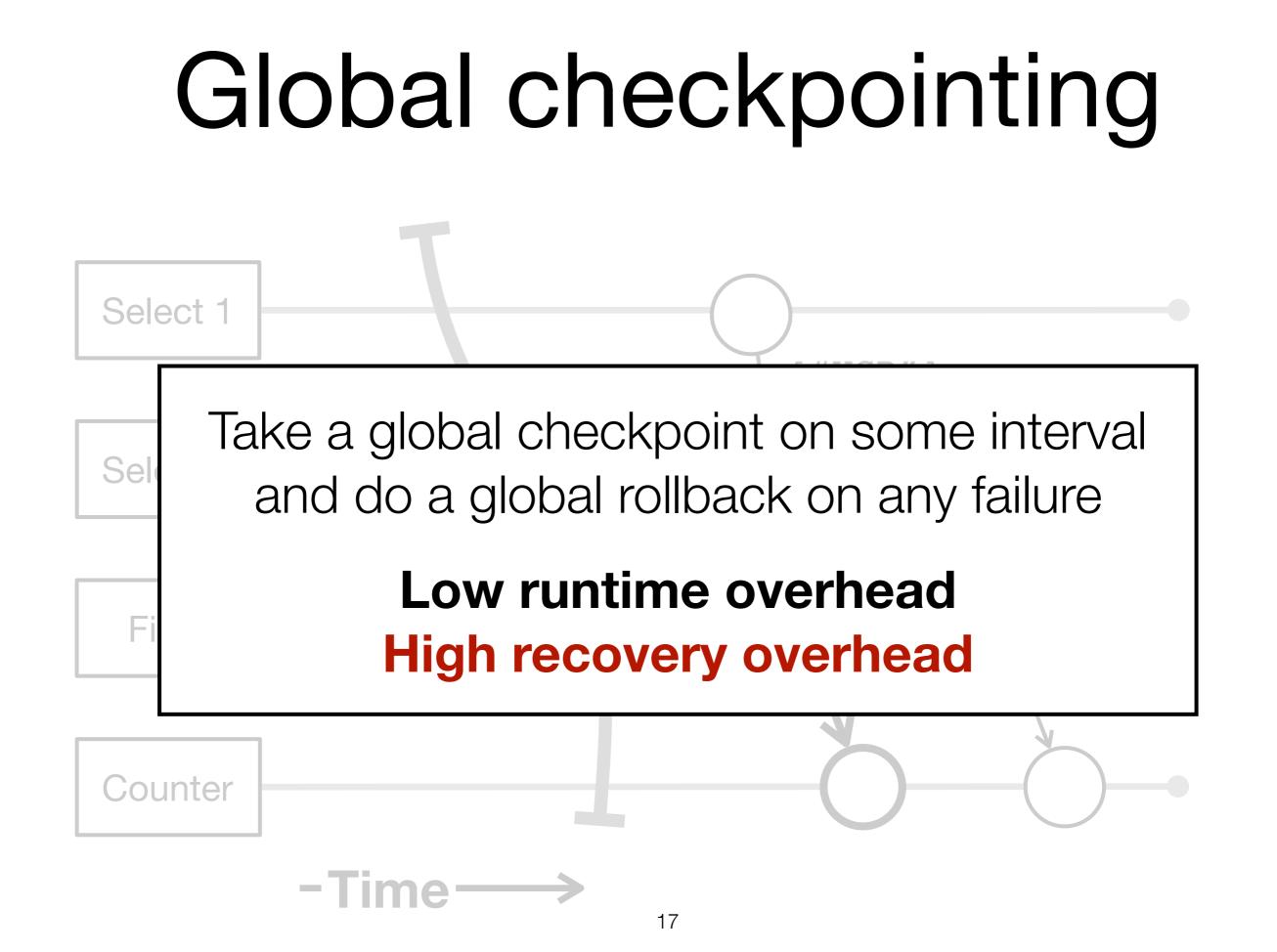




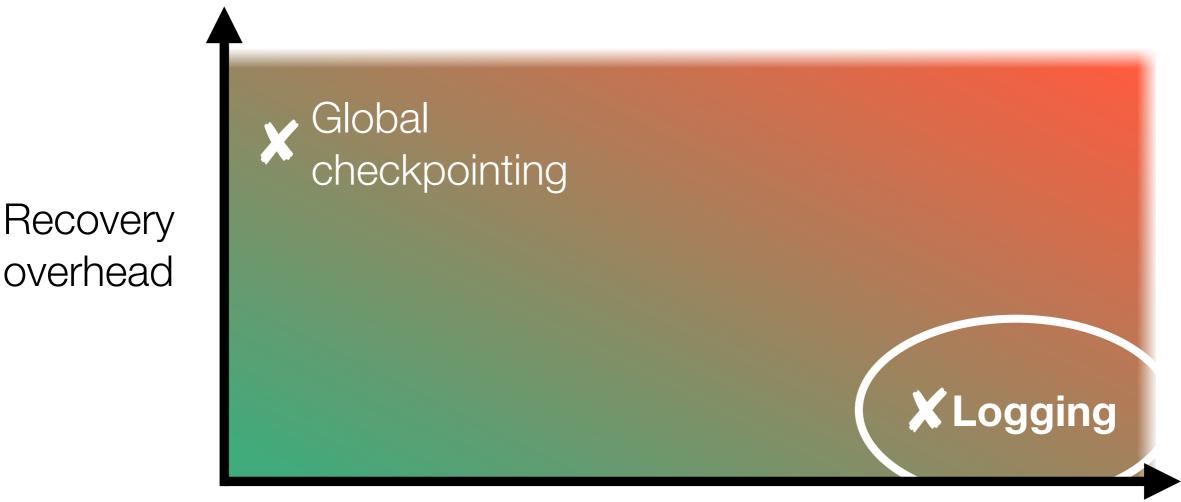




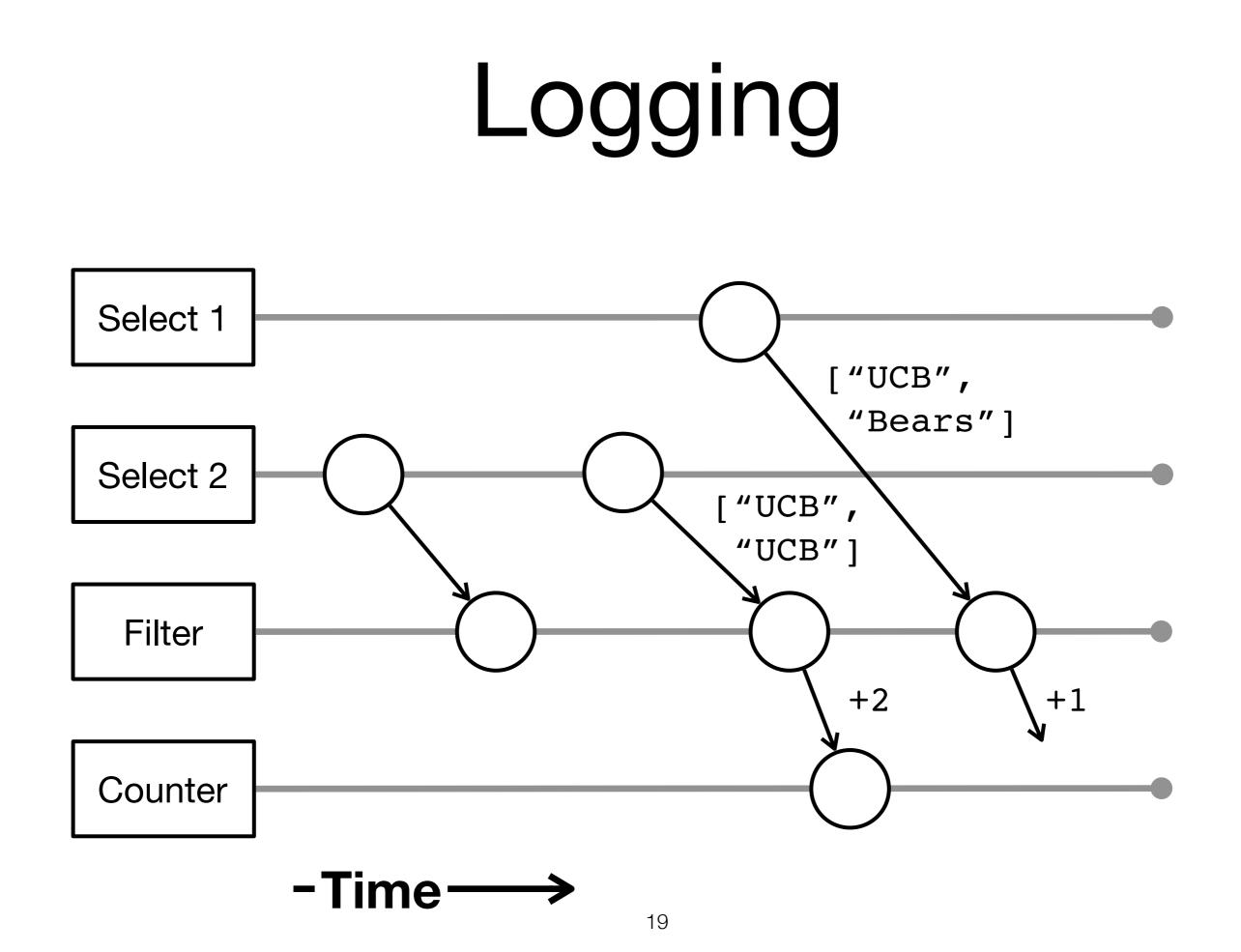


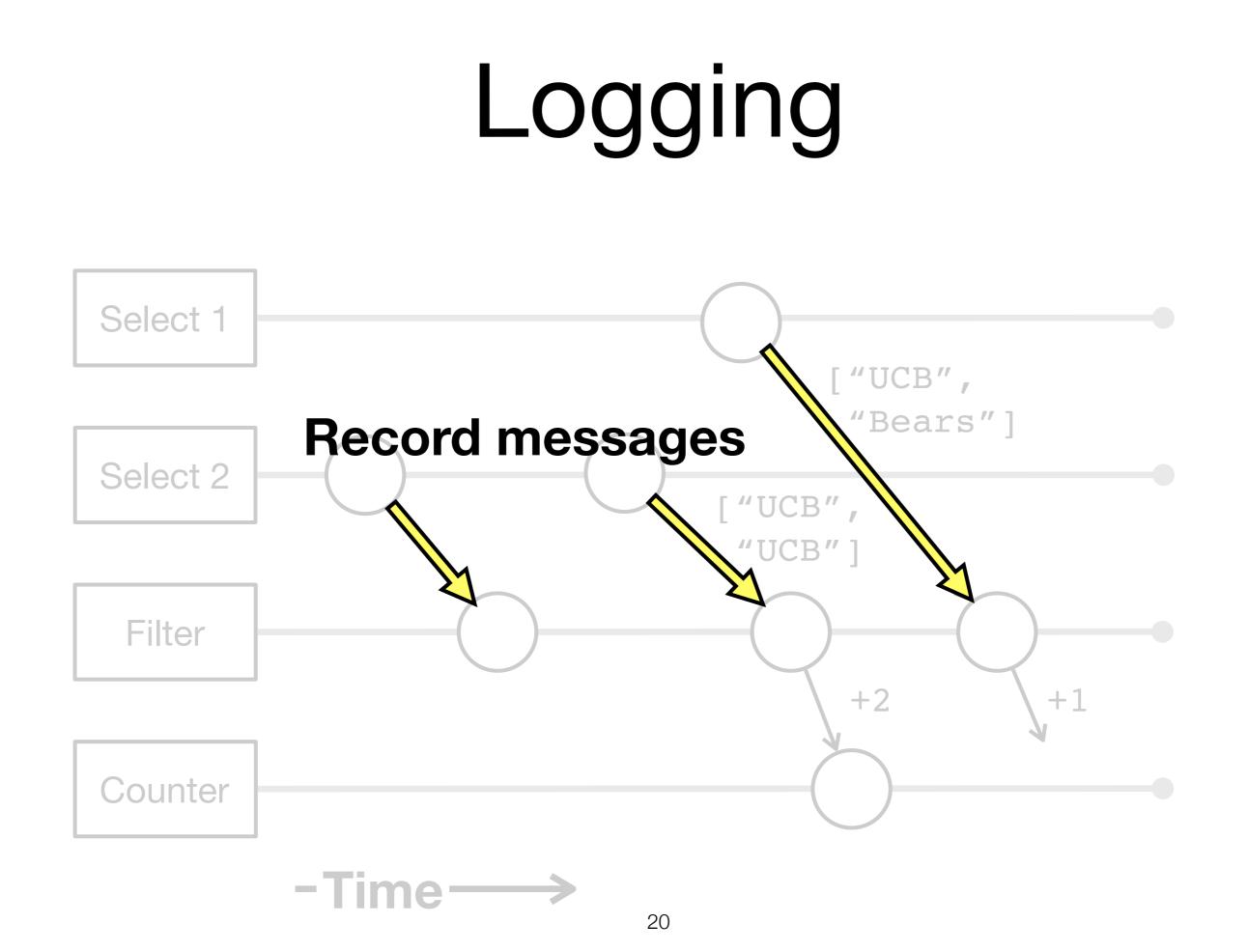


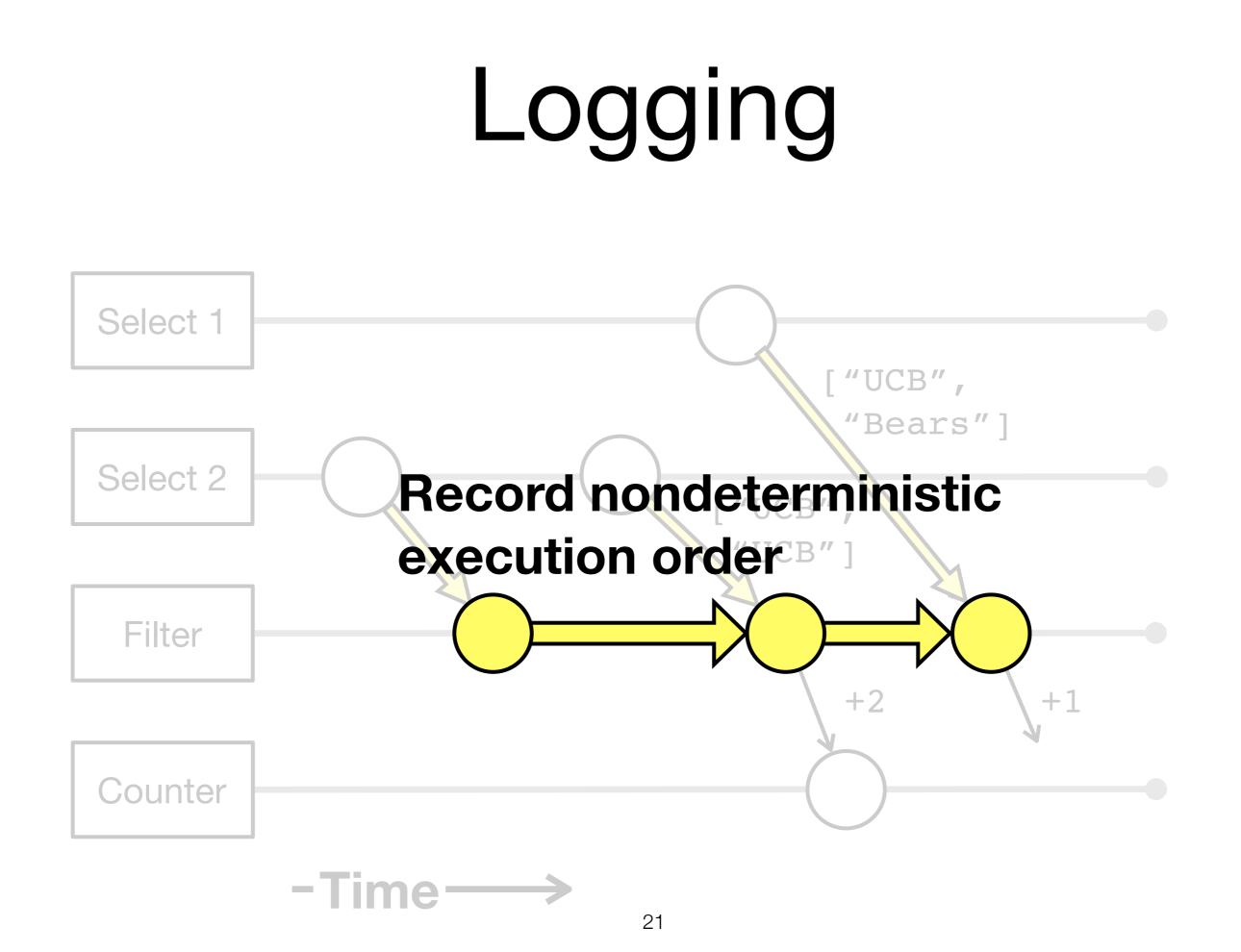
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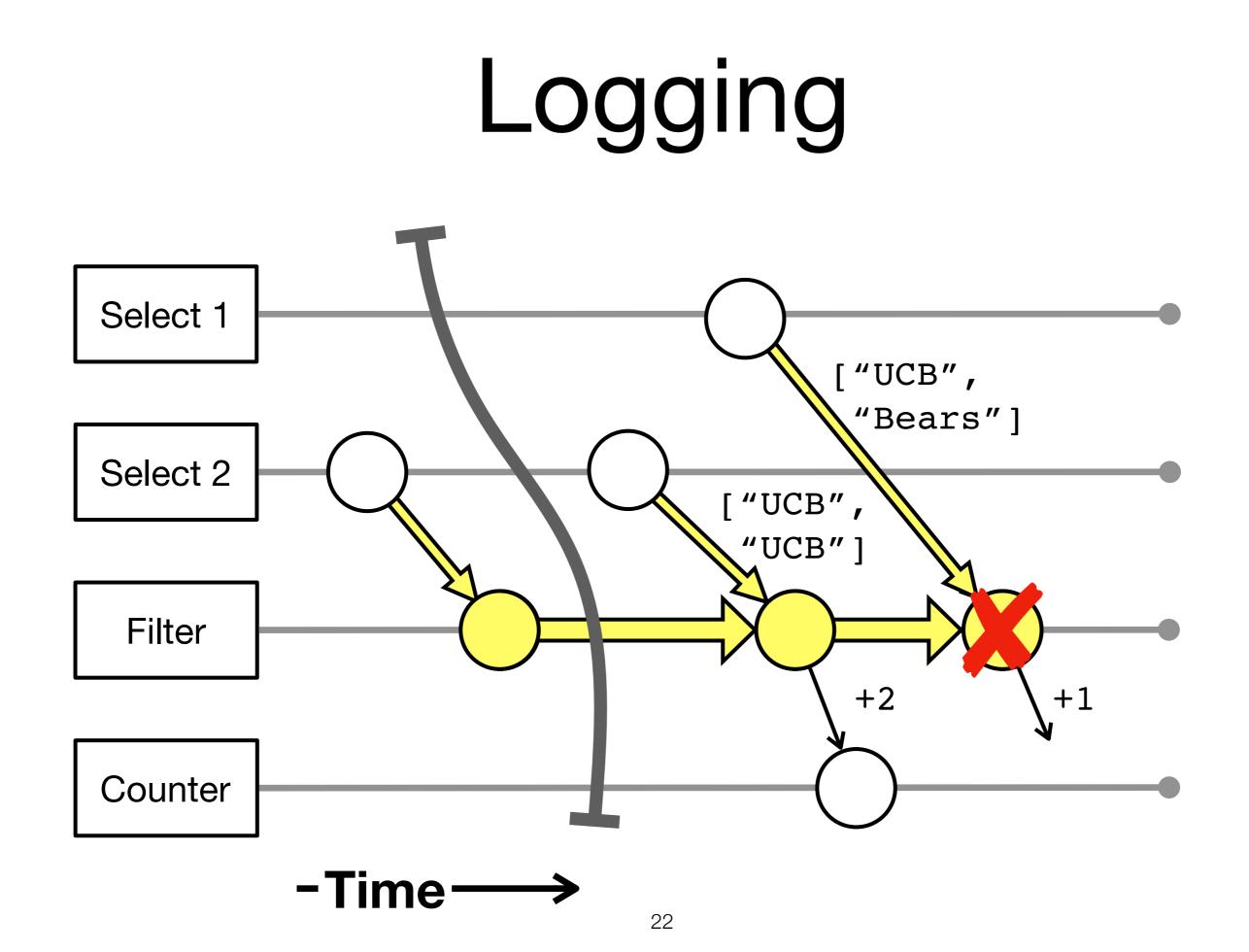


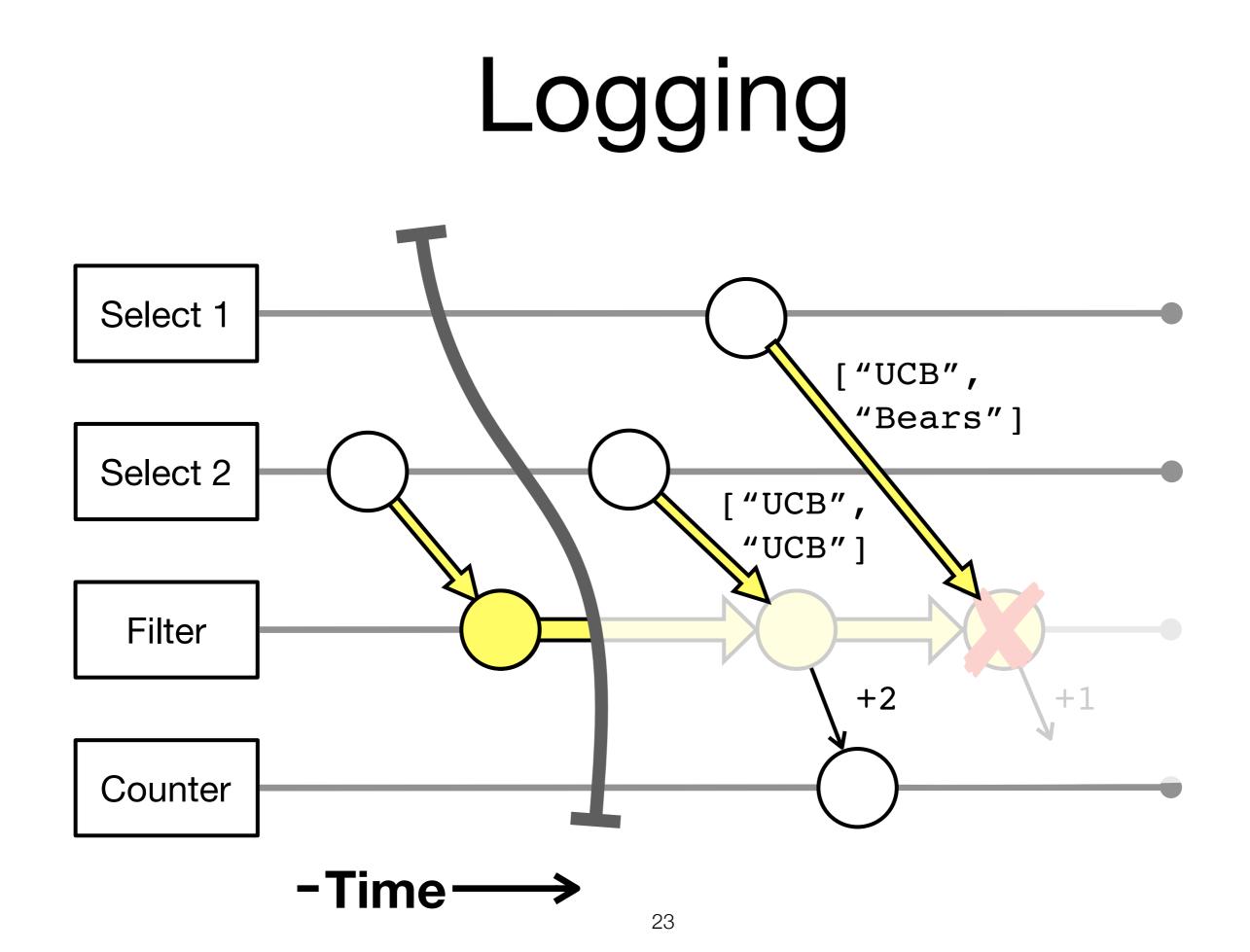
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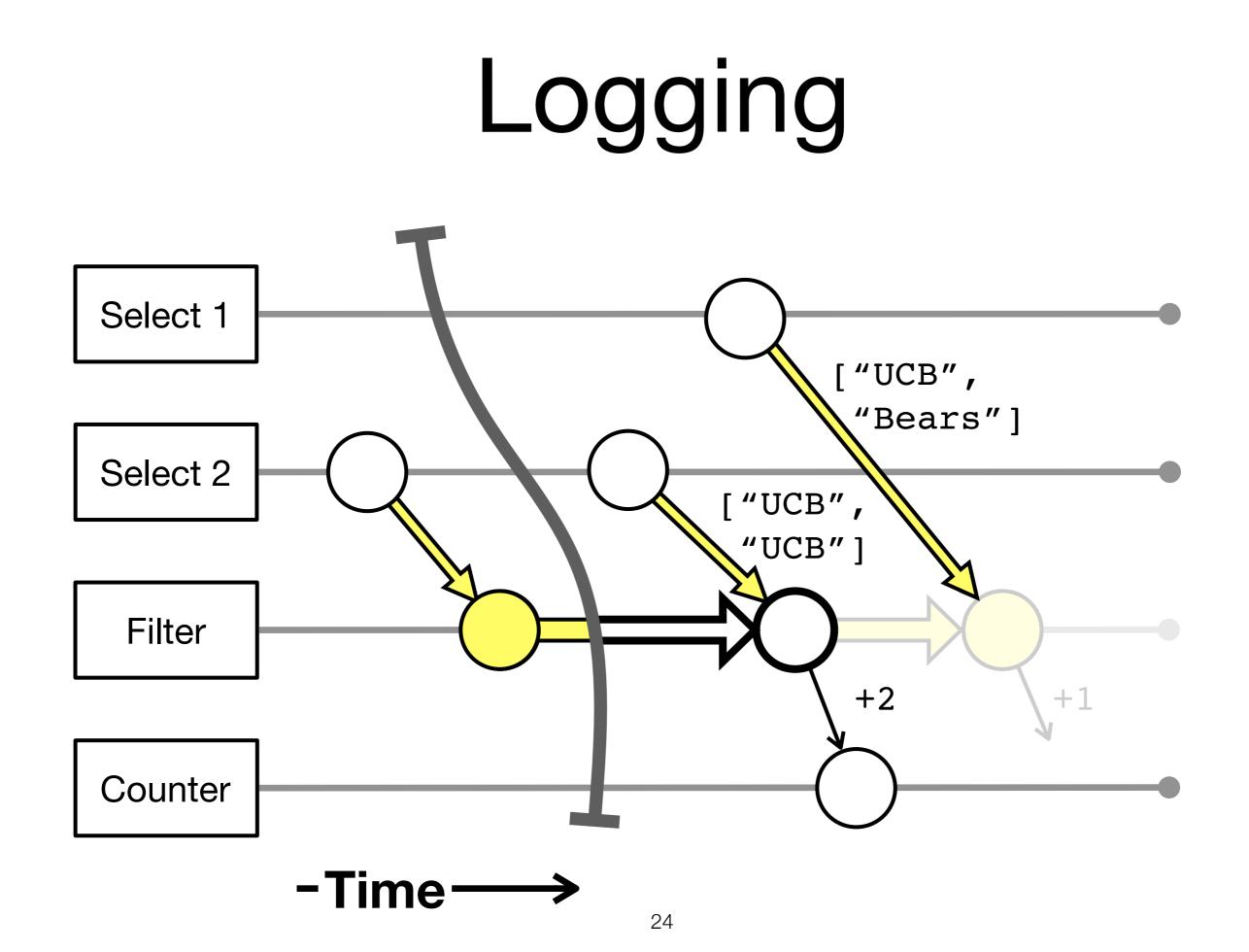


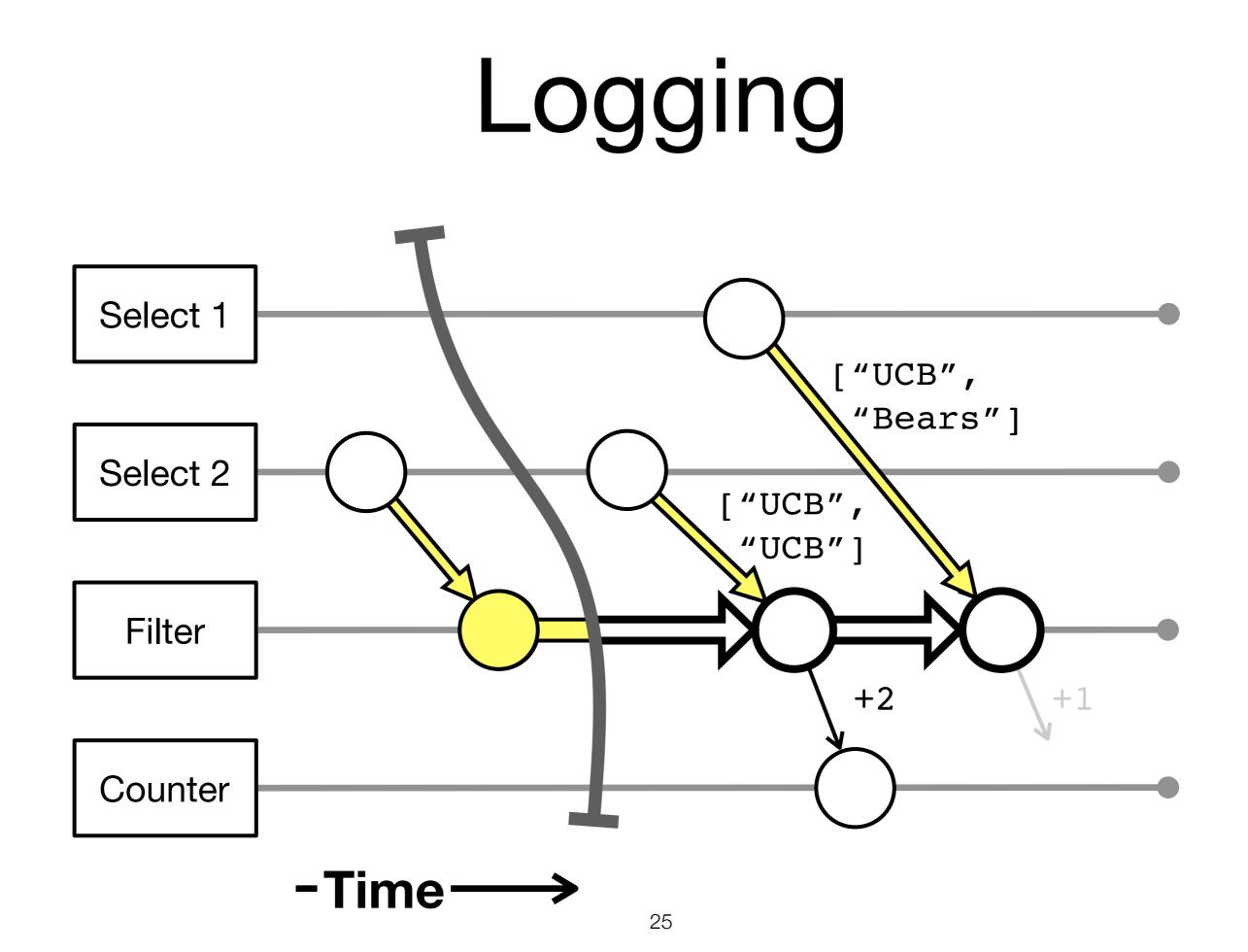


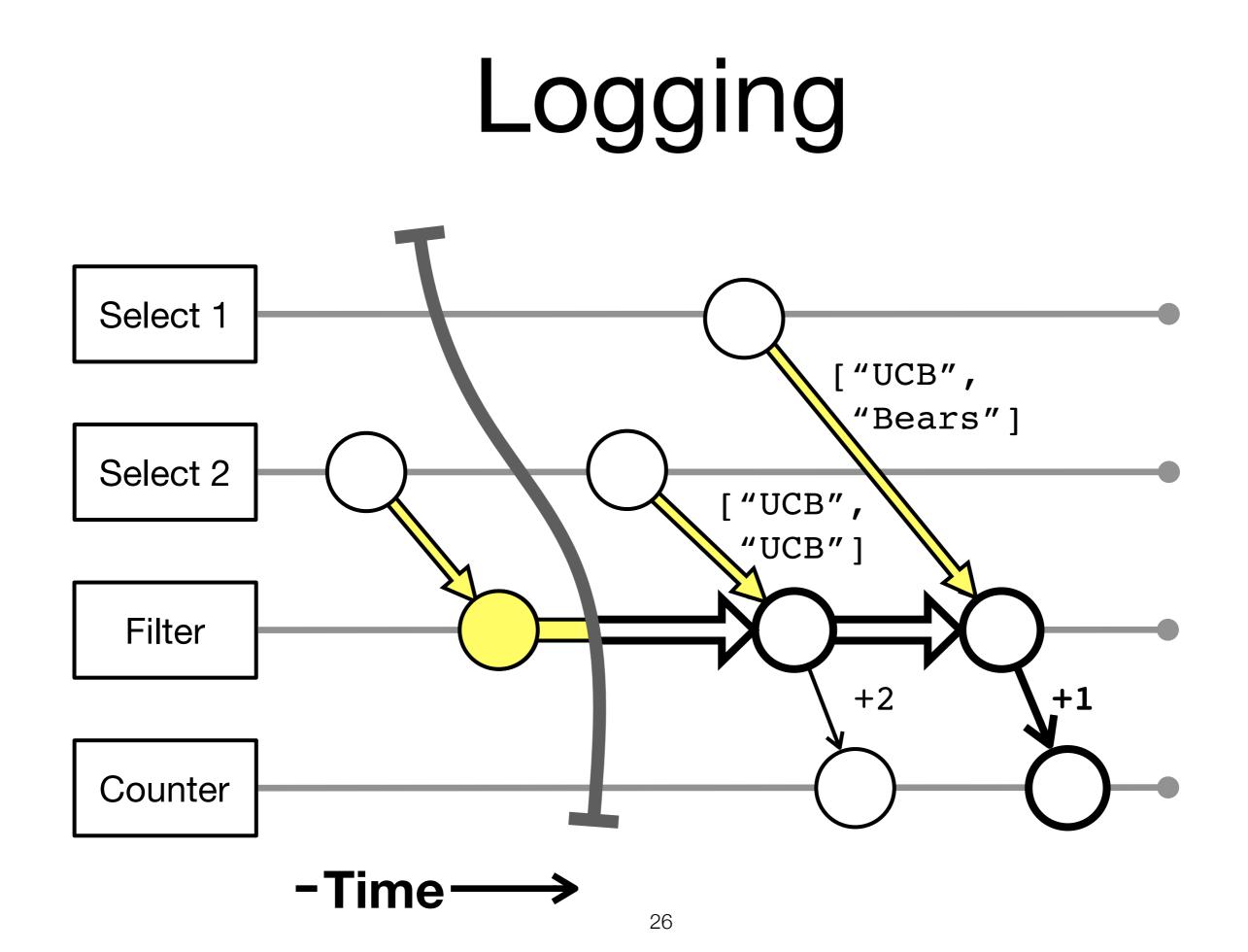


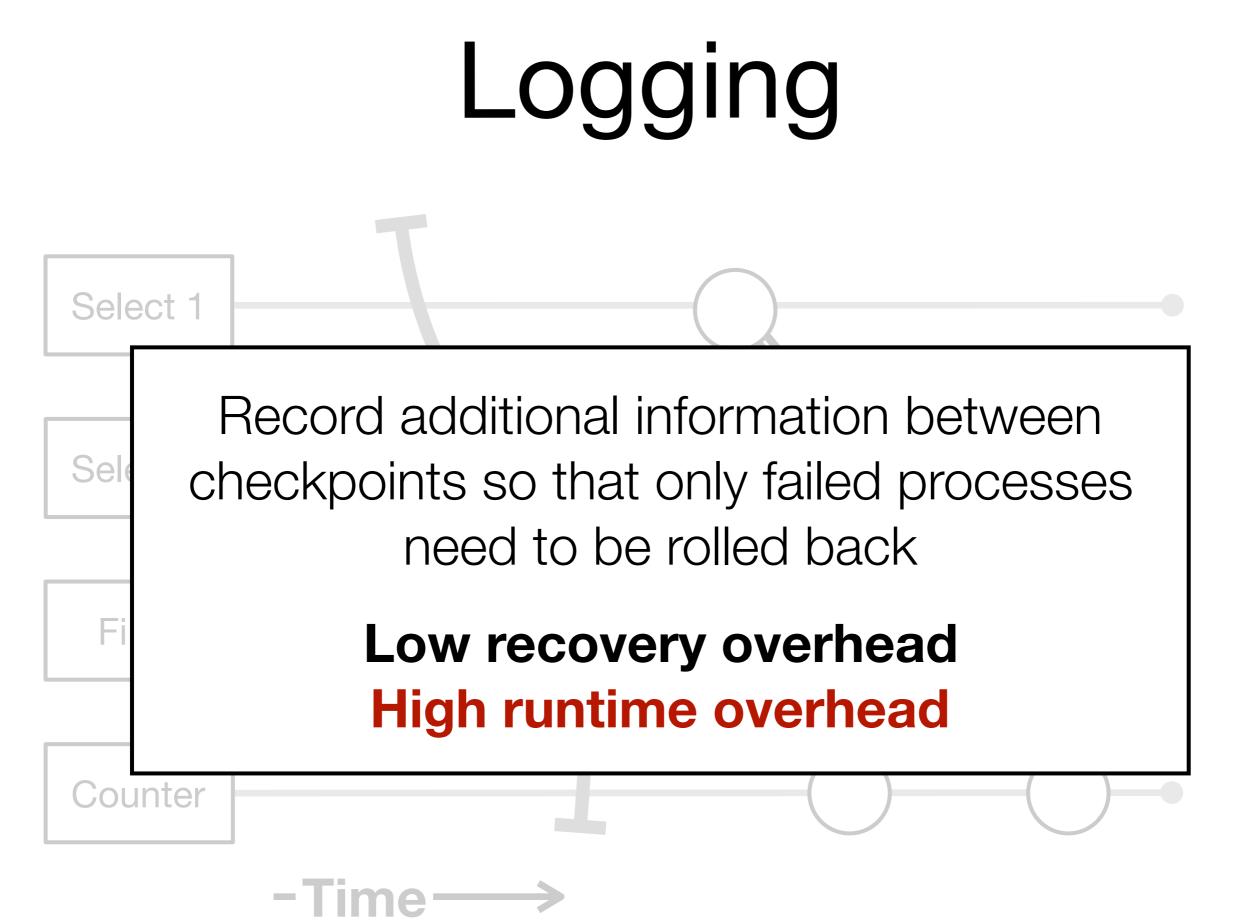


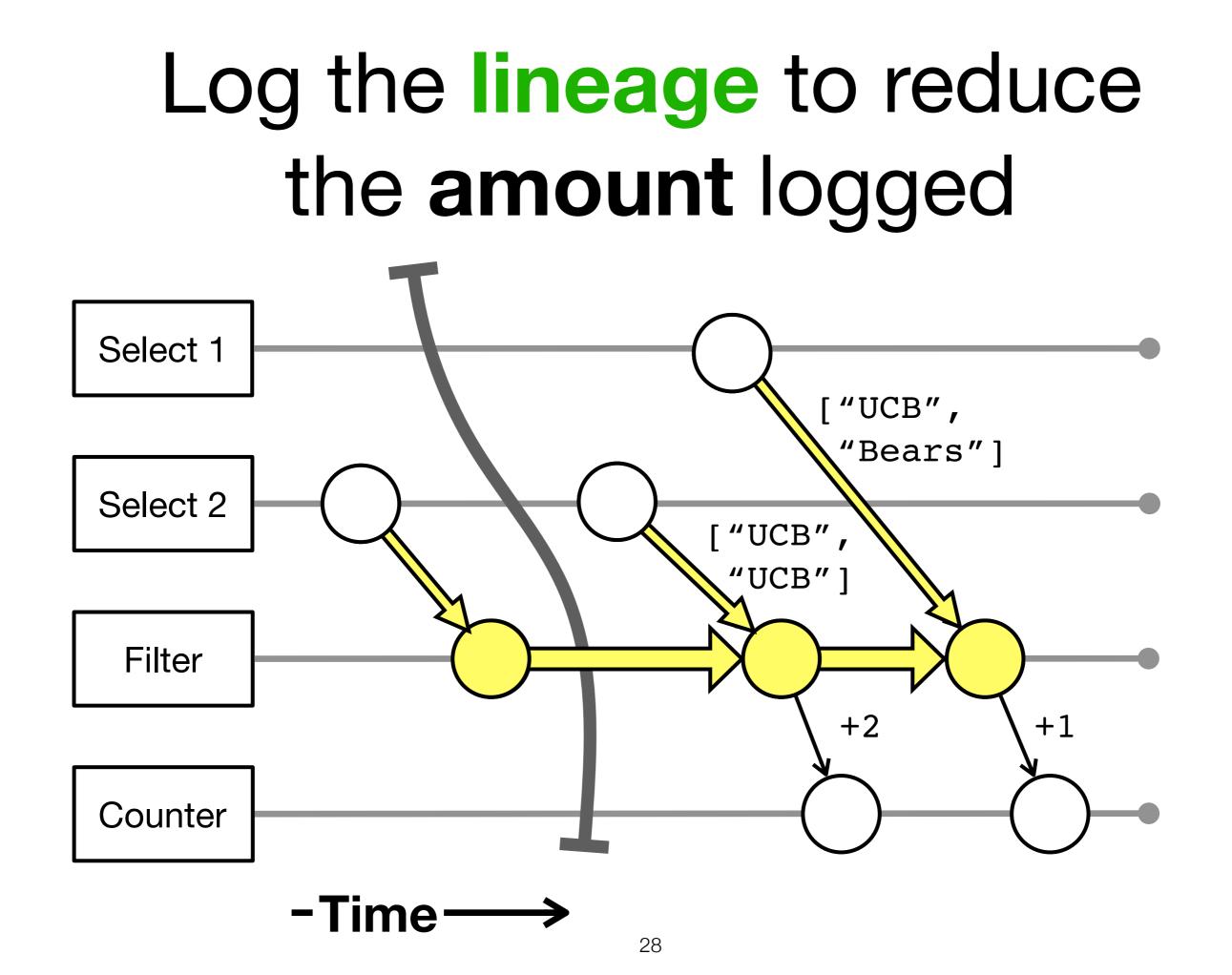


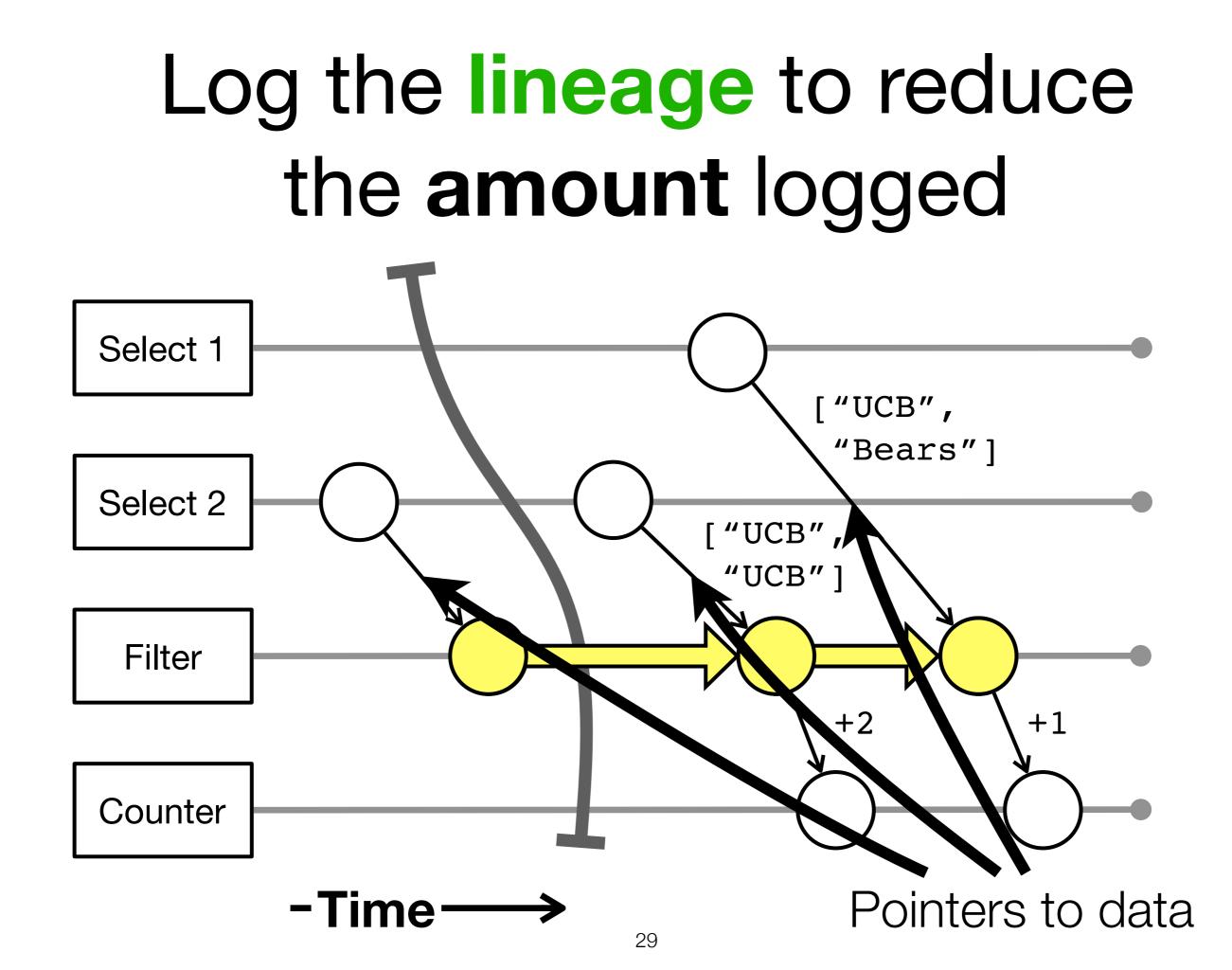




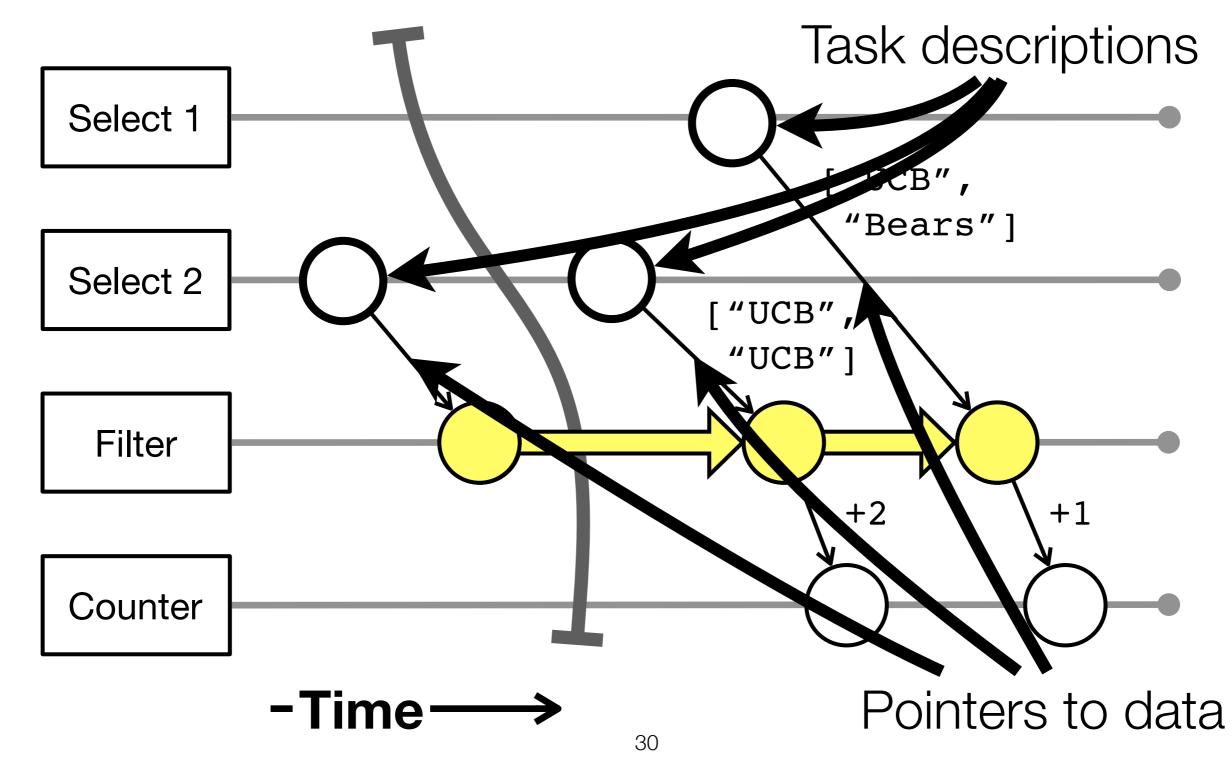


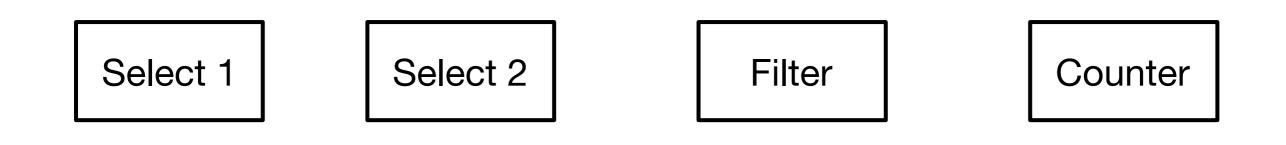


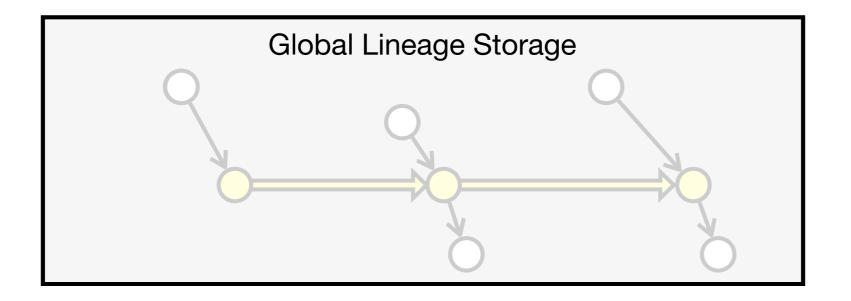


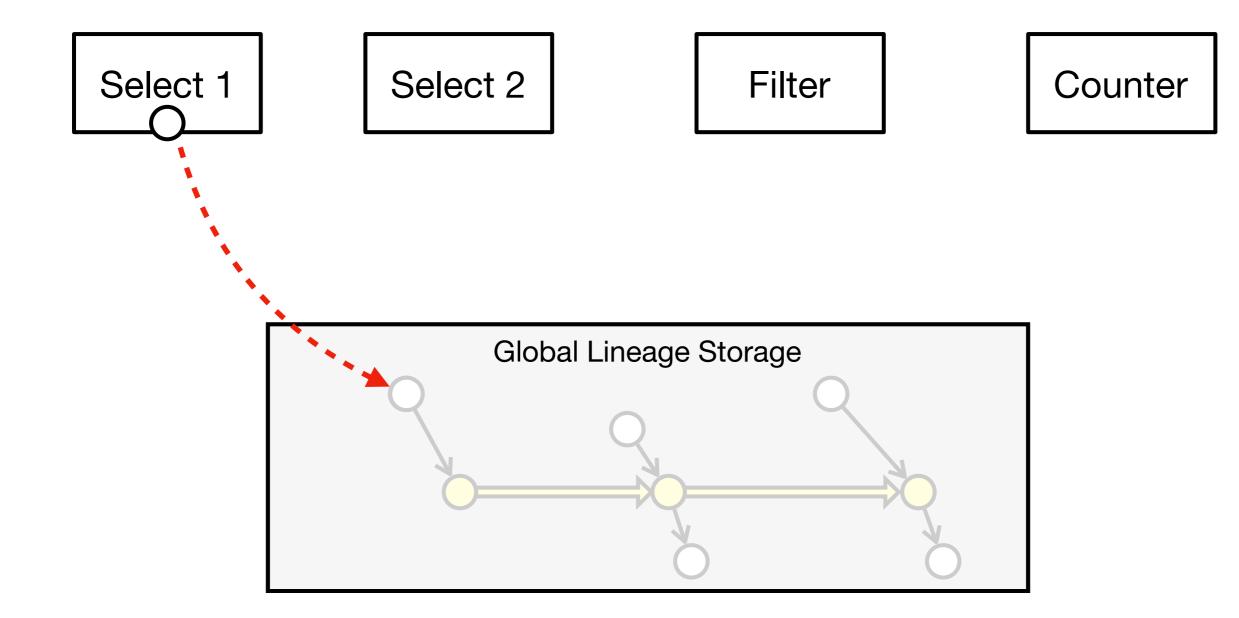


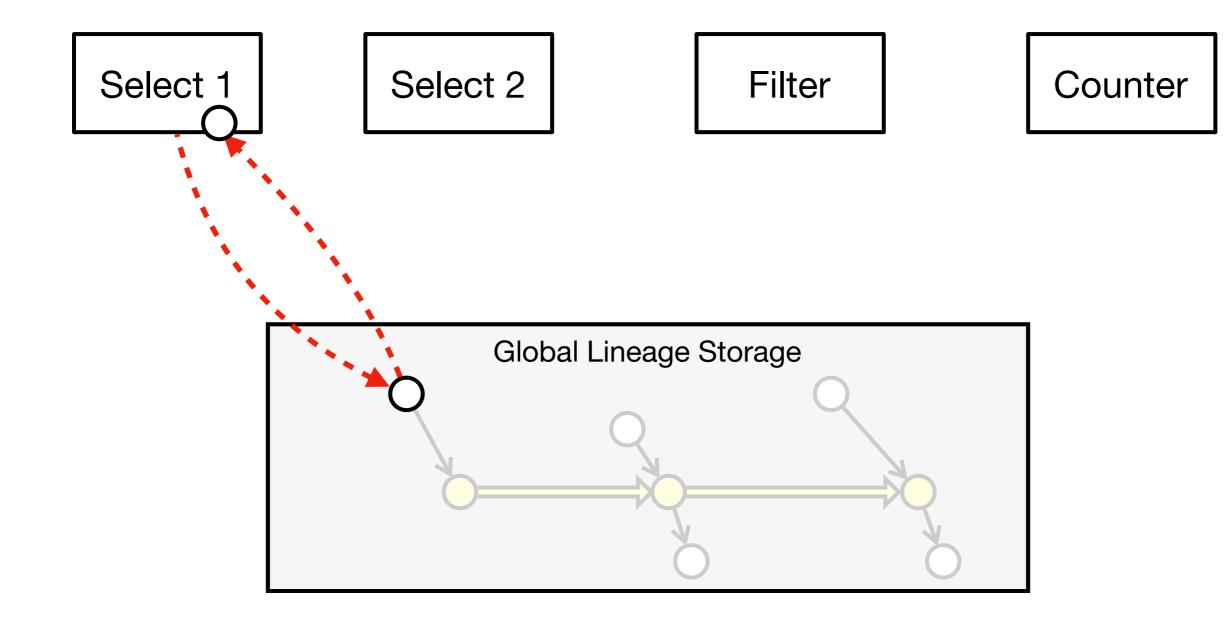
## Log the lineage to reduce the amount logged

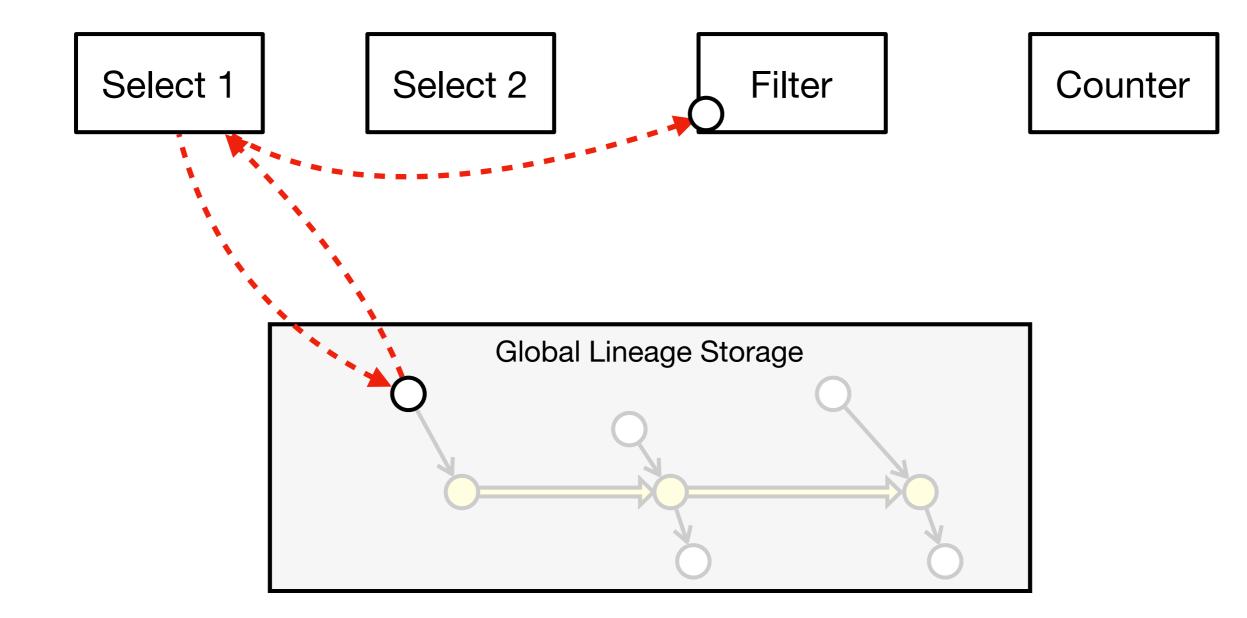


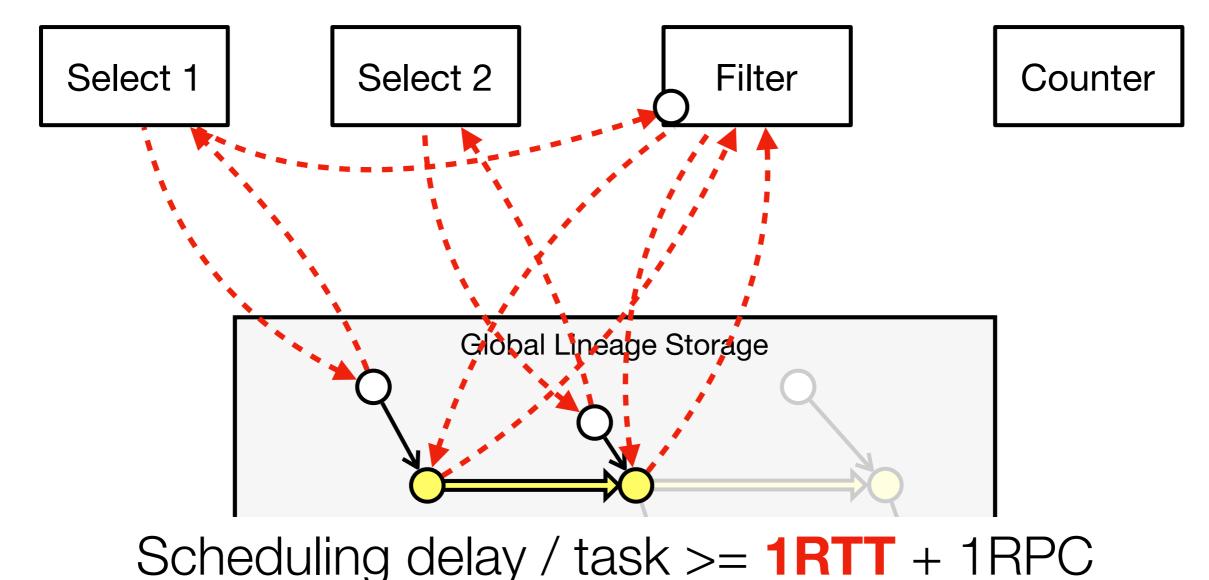












Task latency depends on global storage latency

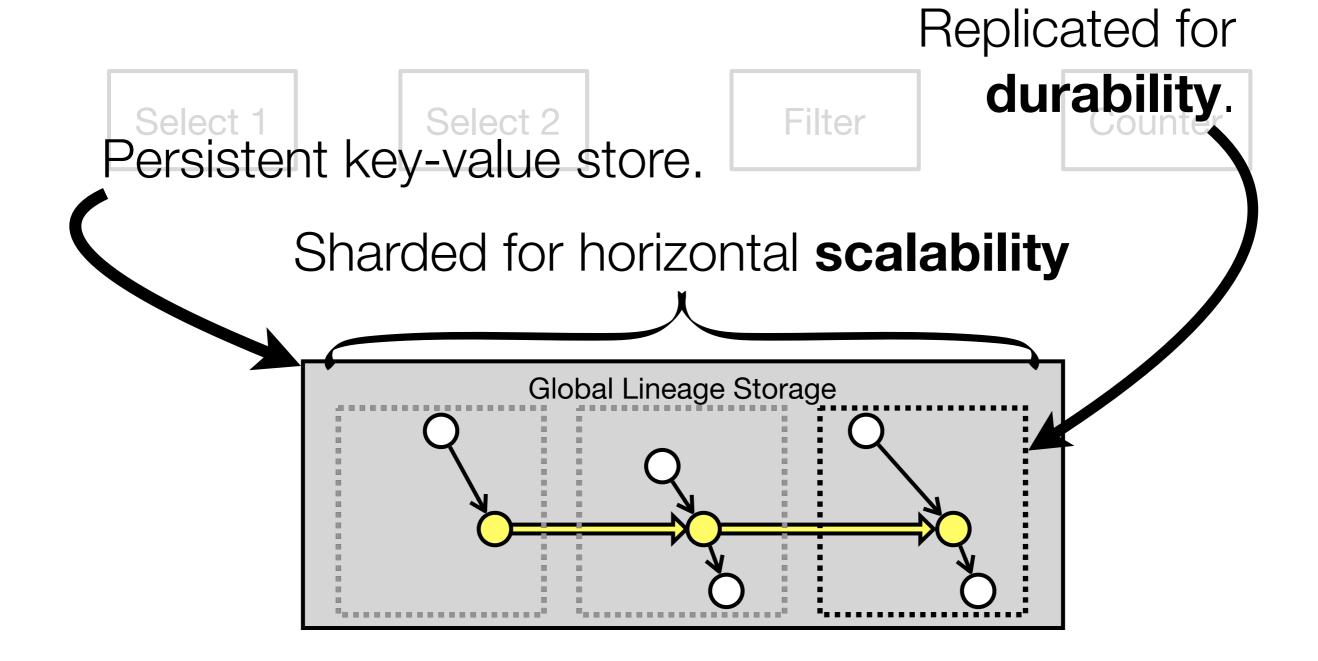
### Lineage stash contribution

How do we achieve both **low runtime** *and* **low recovery overhead** for **fine-grained** data processing applications?

### Solution: Asynchronously log the lineage off the critical path of execution.

Lineage reconstruction to reduce amount logged

Causal logging to log nondeterminism



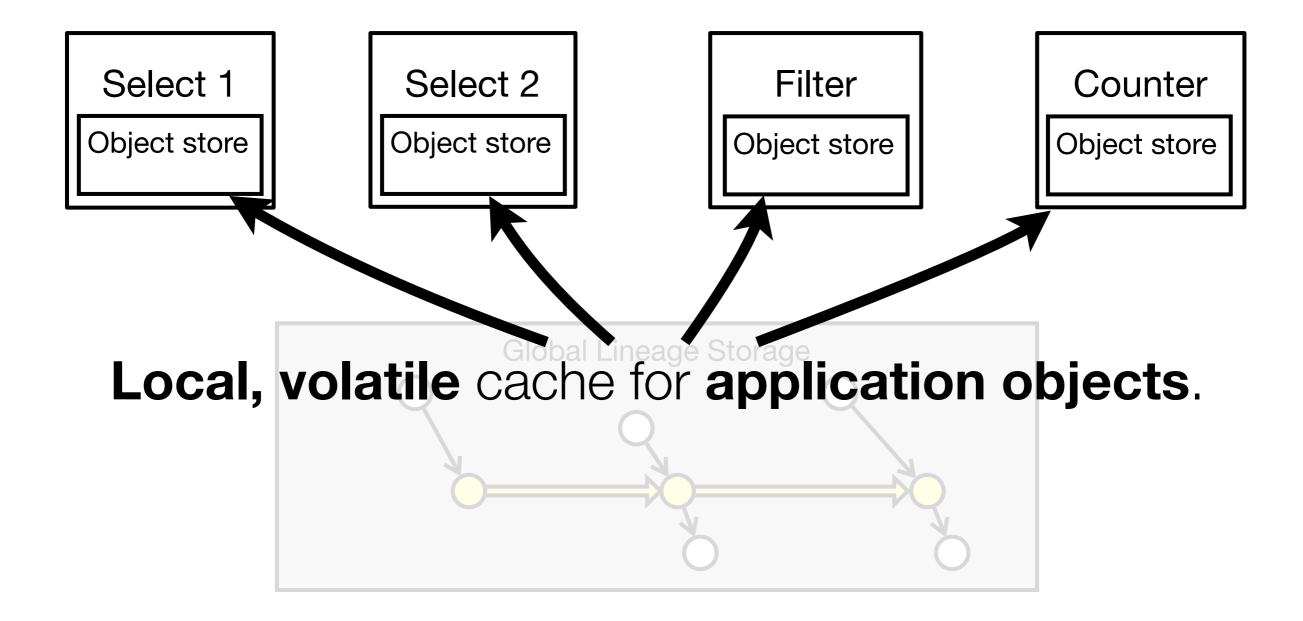


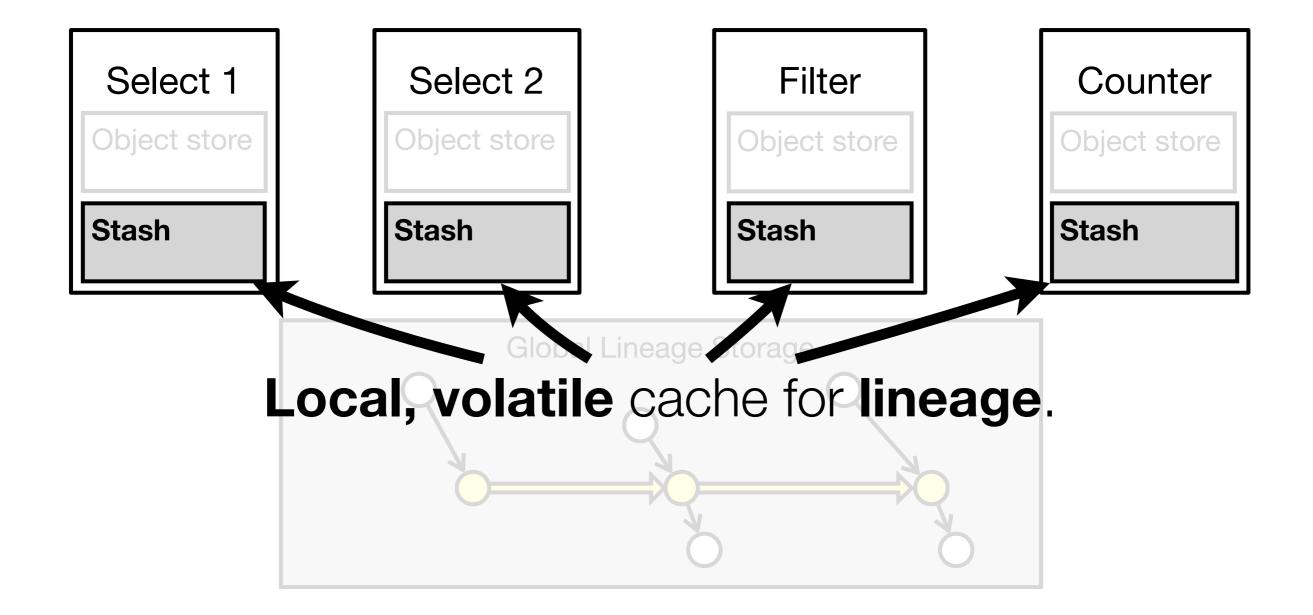


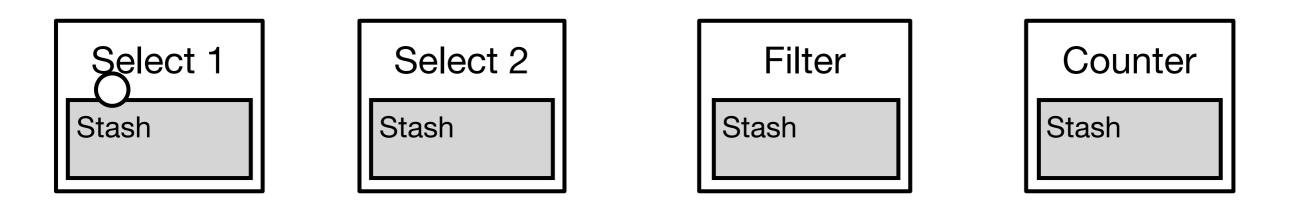






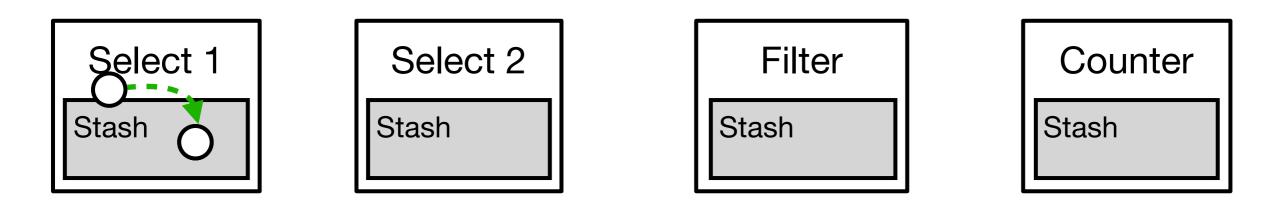






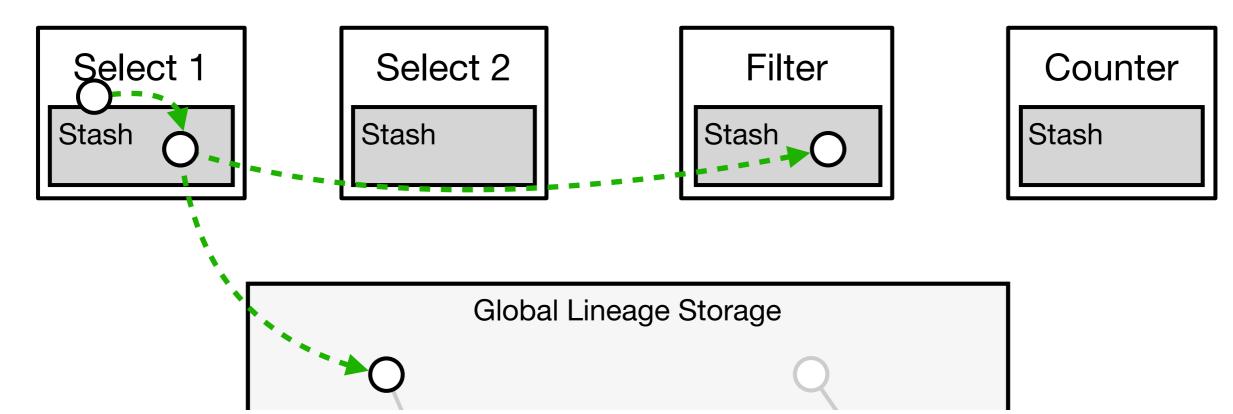
Global Lineage Storage

 (1) Write lineage to local, volatile lineage stash.
(2) Asynchronously flush to remote storage. Scheduling delay / task = 1RTT + 1RPC
Task latency independent of global storage latency



Global Lineage Storage

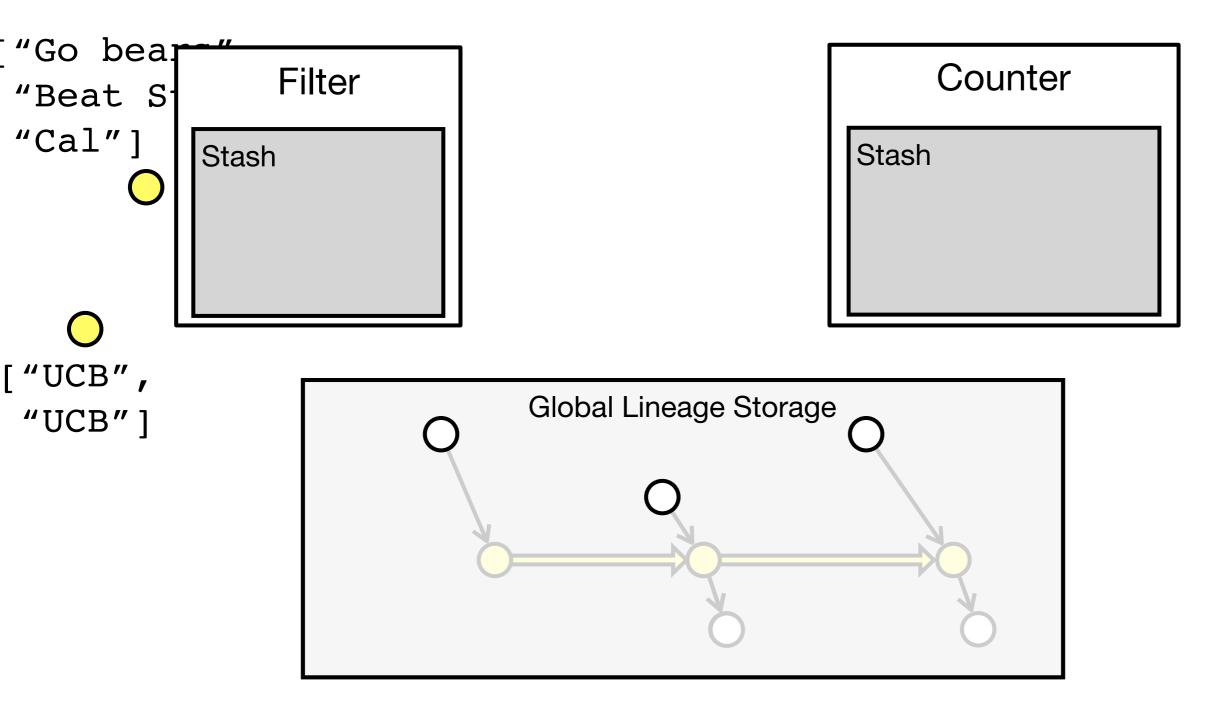
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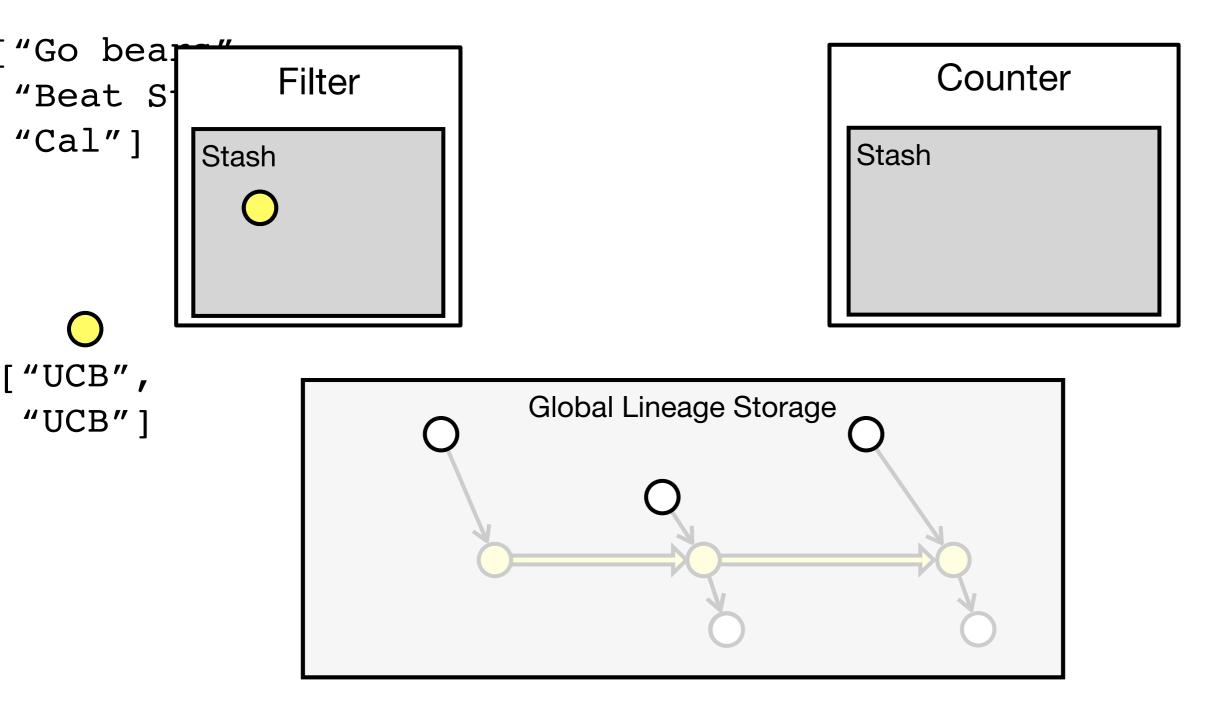


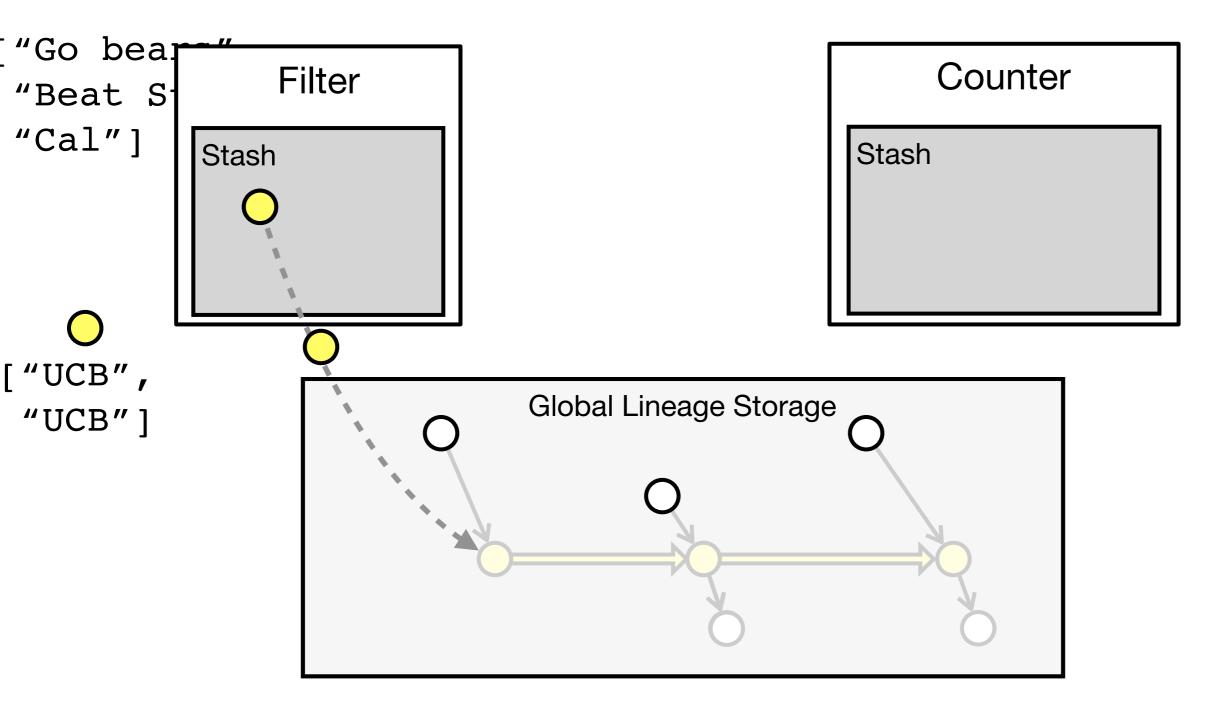
(1) Write lineage to **local, volatile** lineage stash.

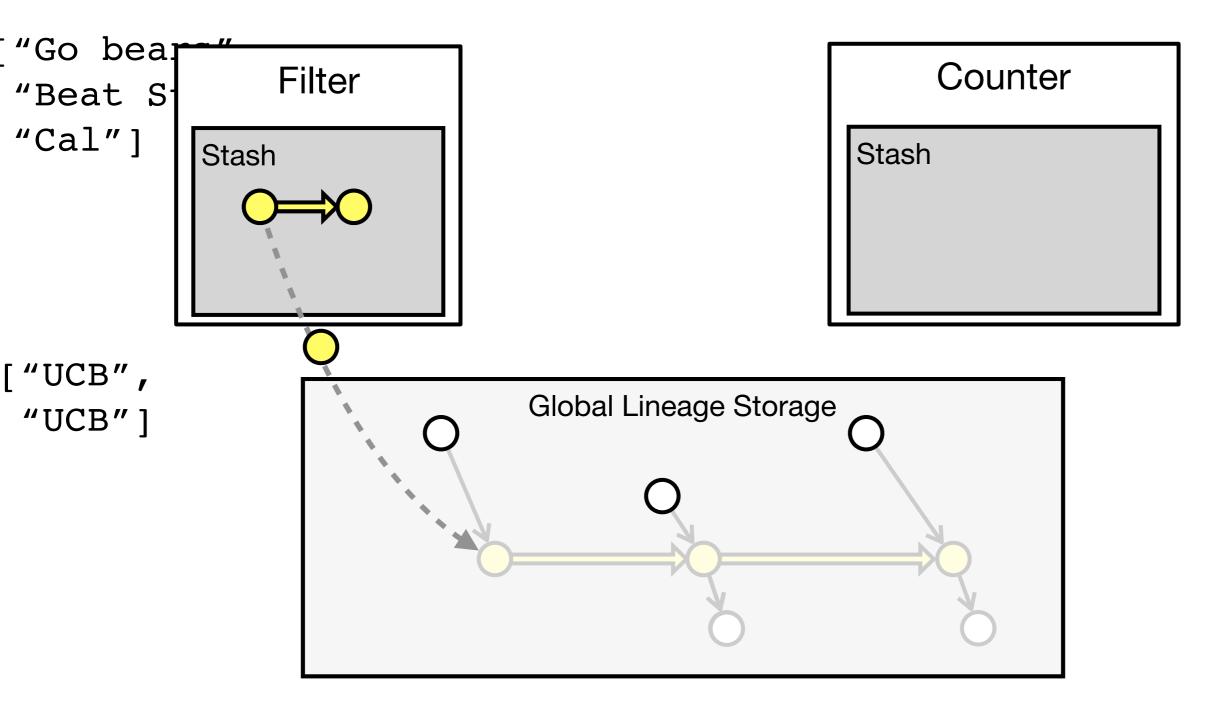
(2) **Asynchronously** flush to **remote** storage.

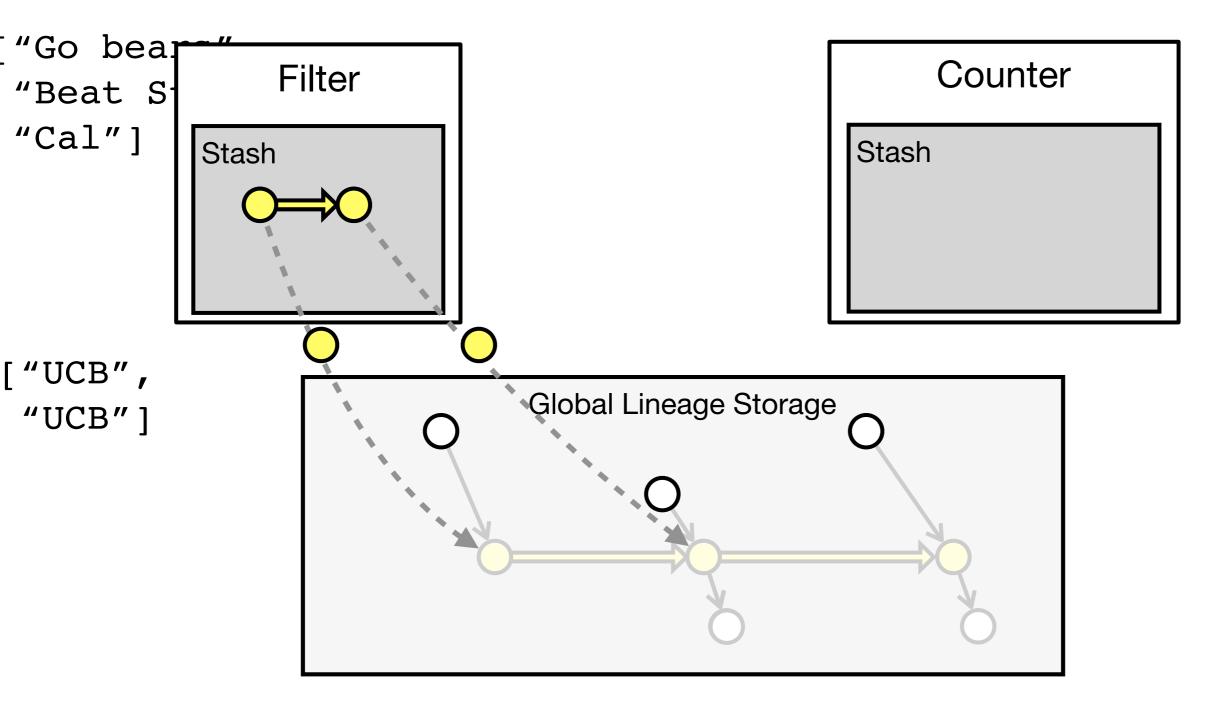
Scheduling delay / task = <del>1RTT</del> + **1RPC** Task latency **independent of** global storage latency

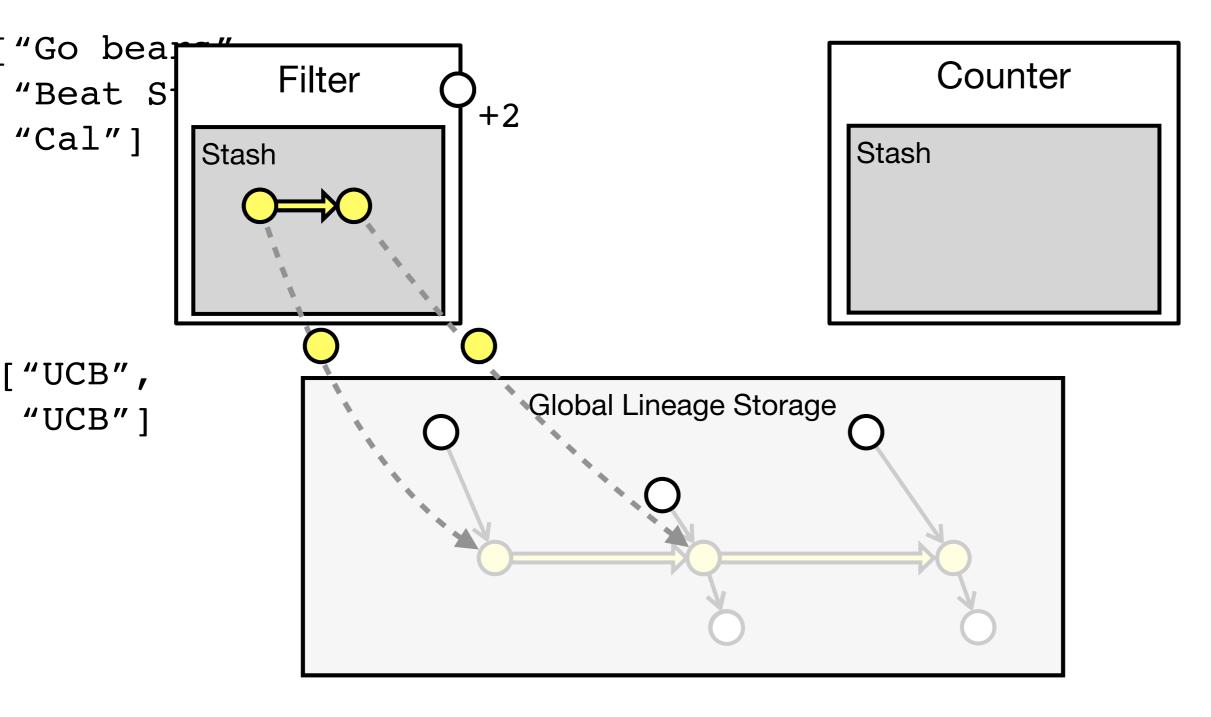


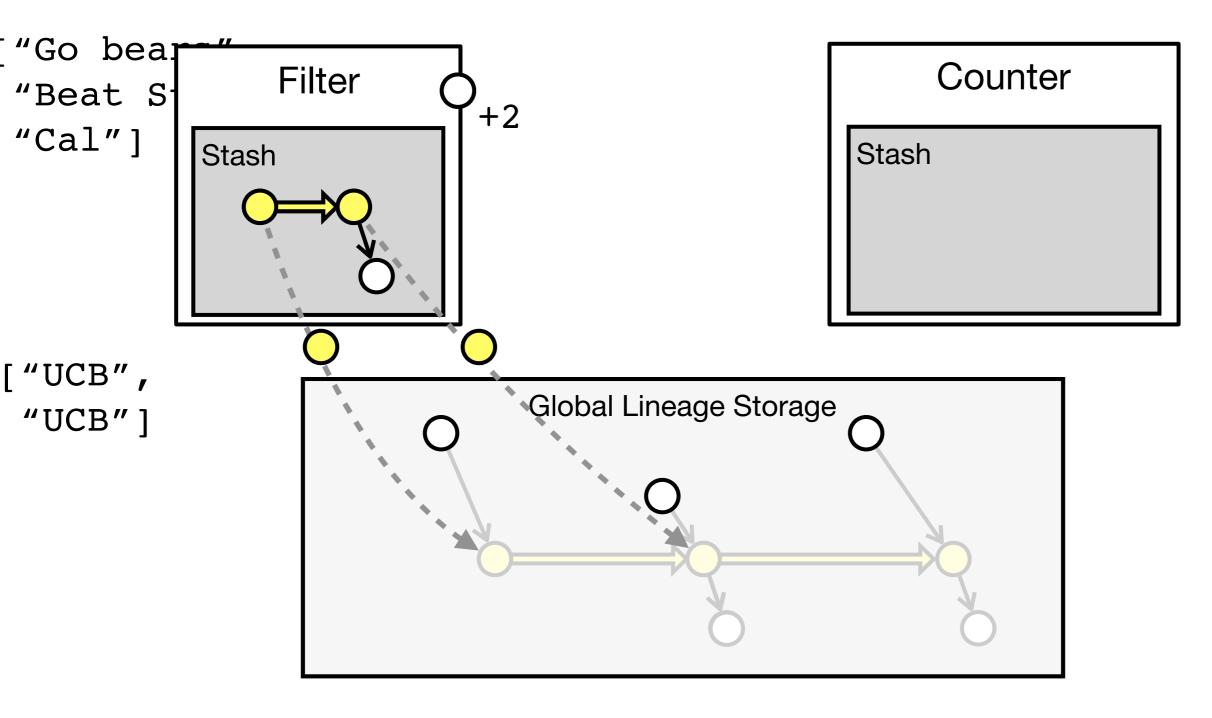


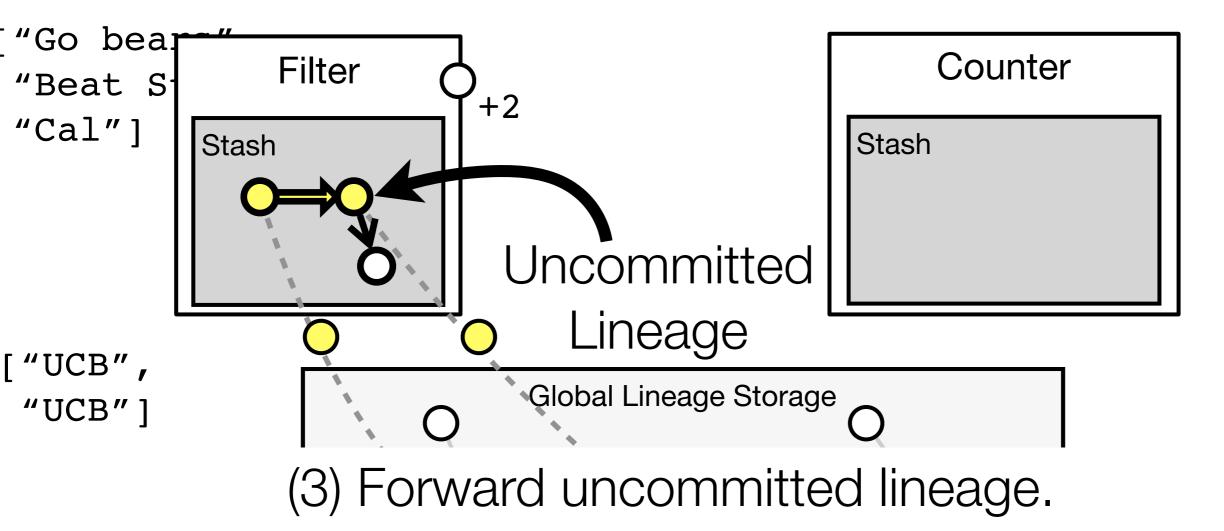


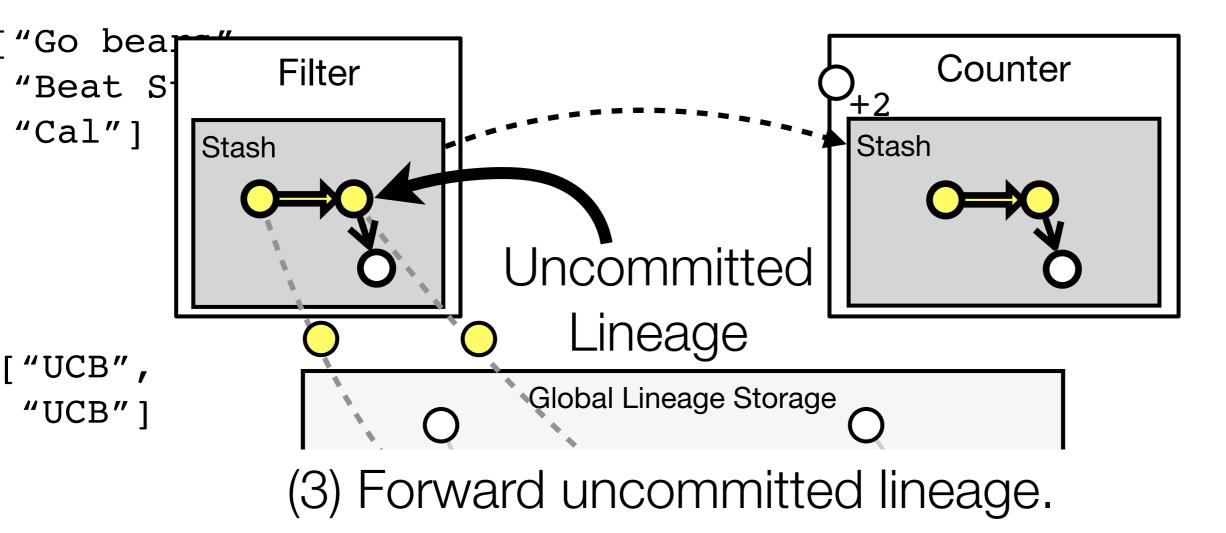


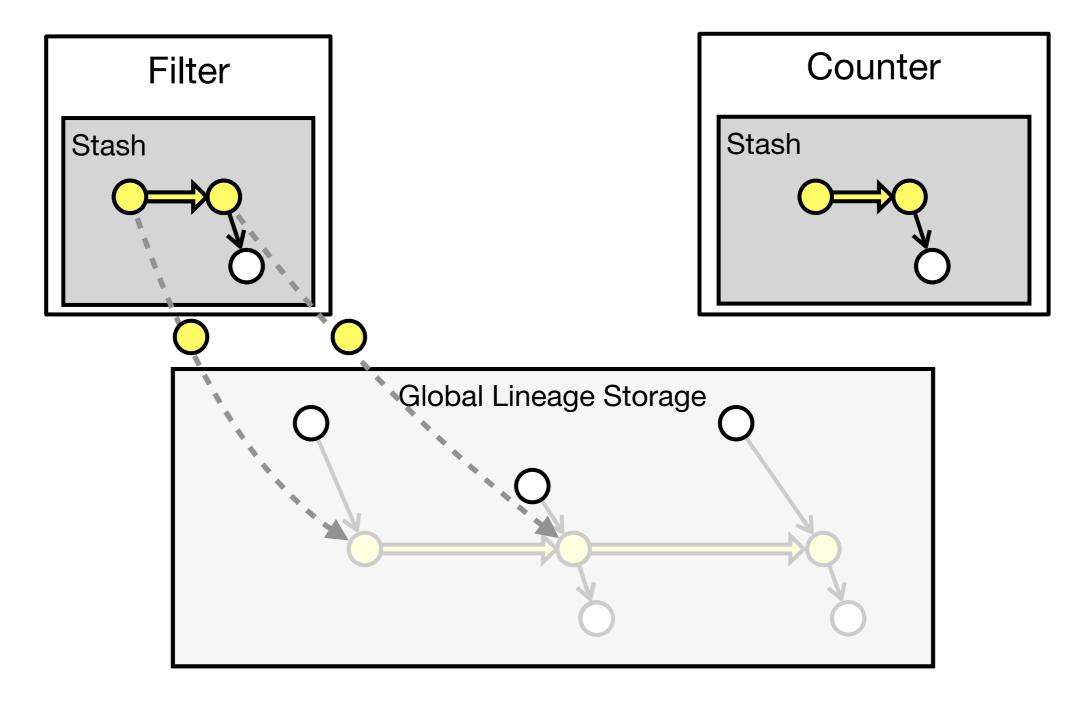


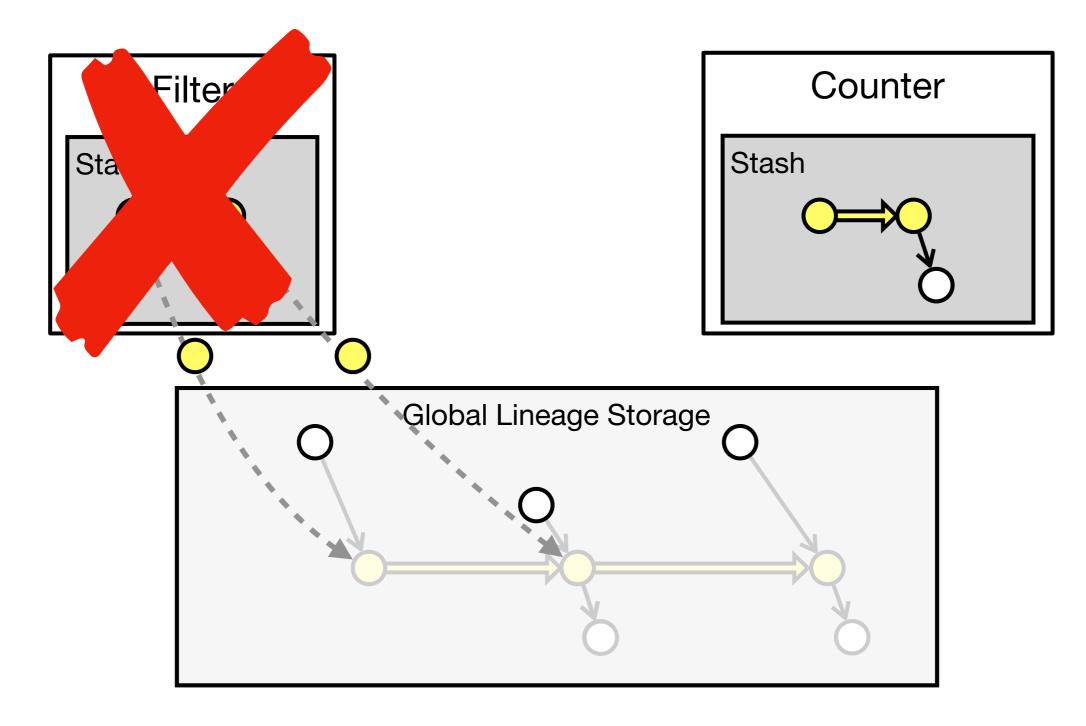


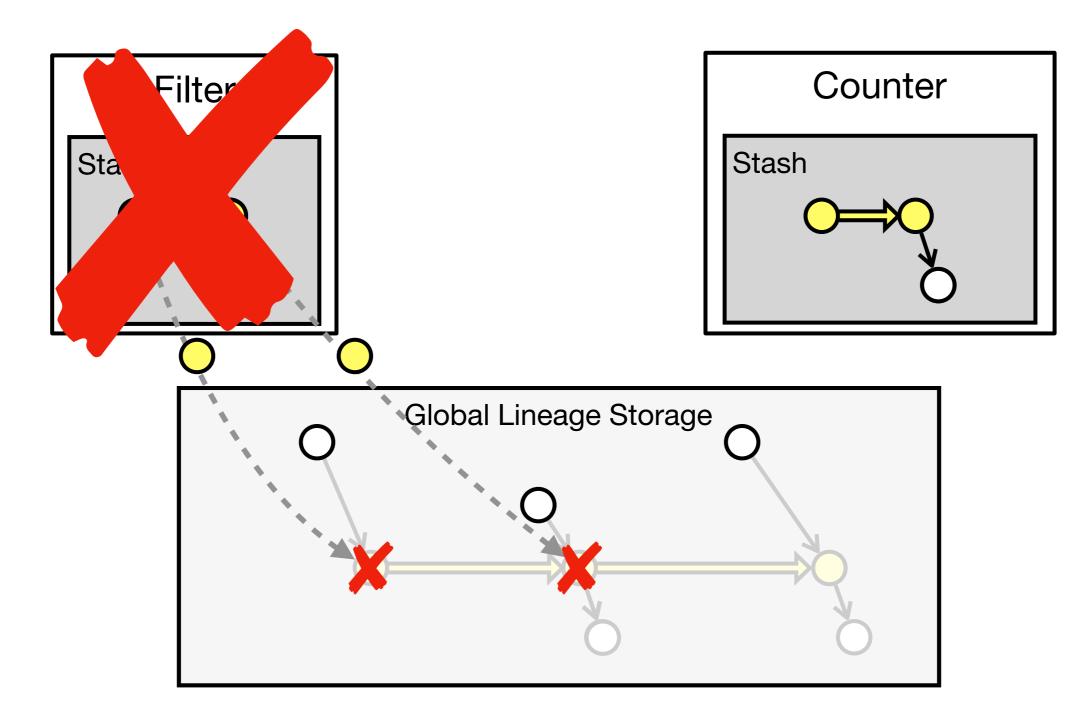


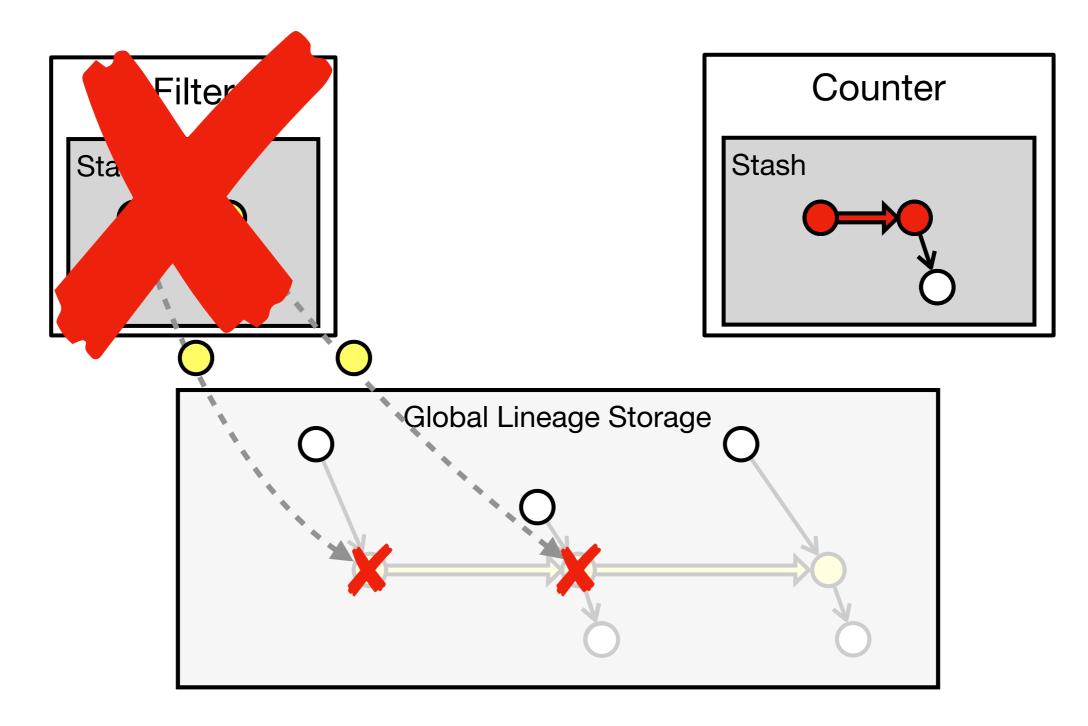


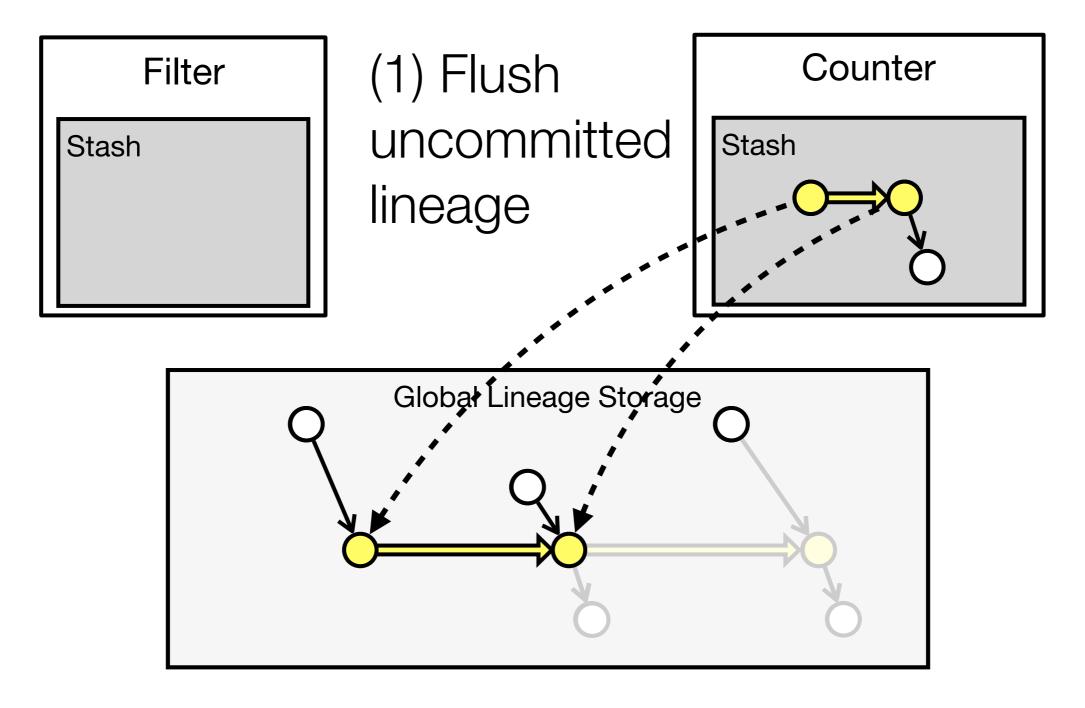


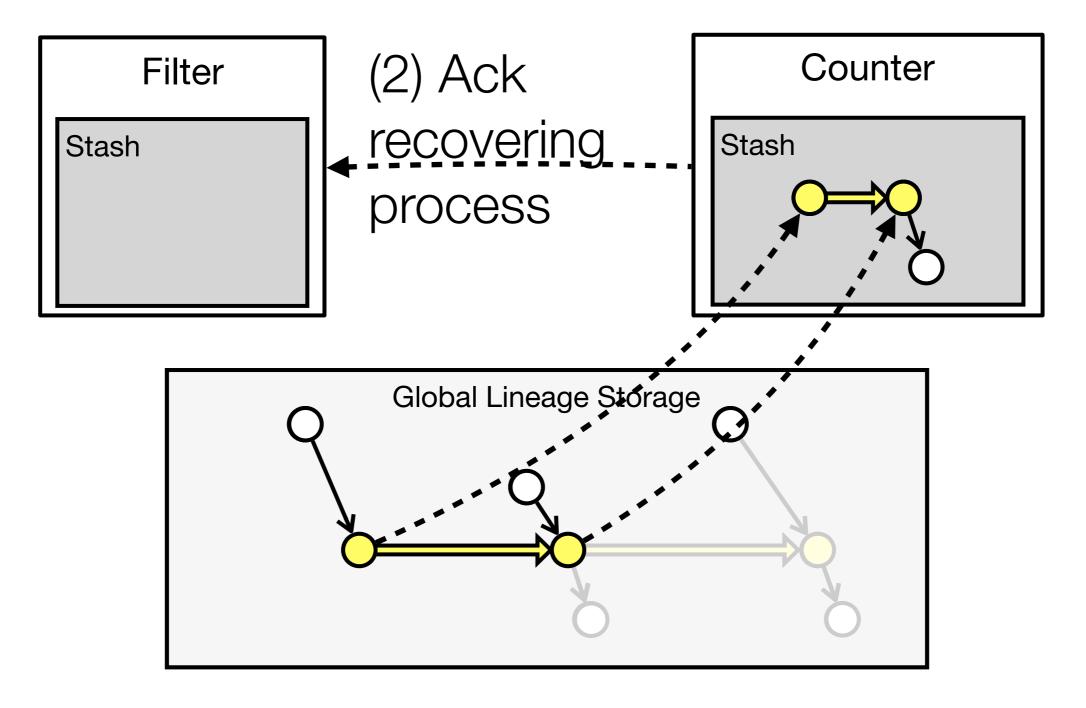


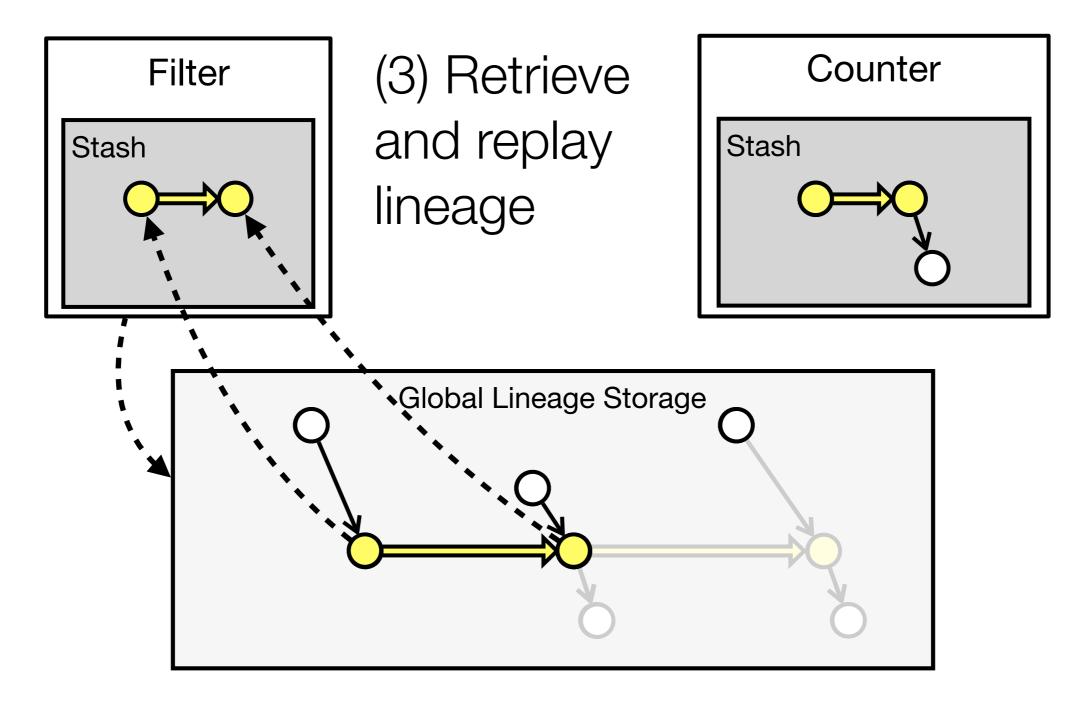




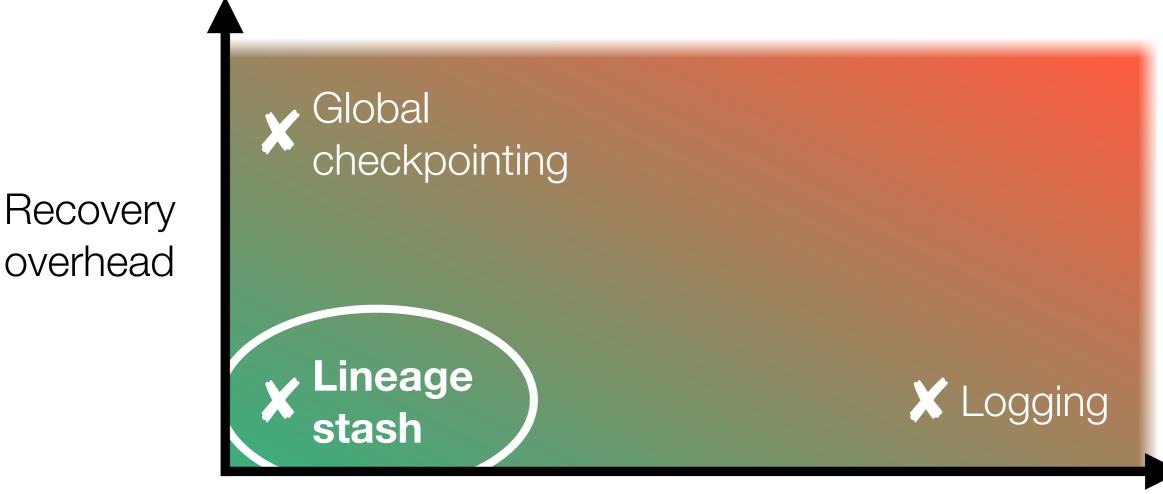




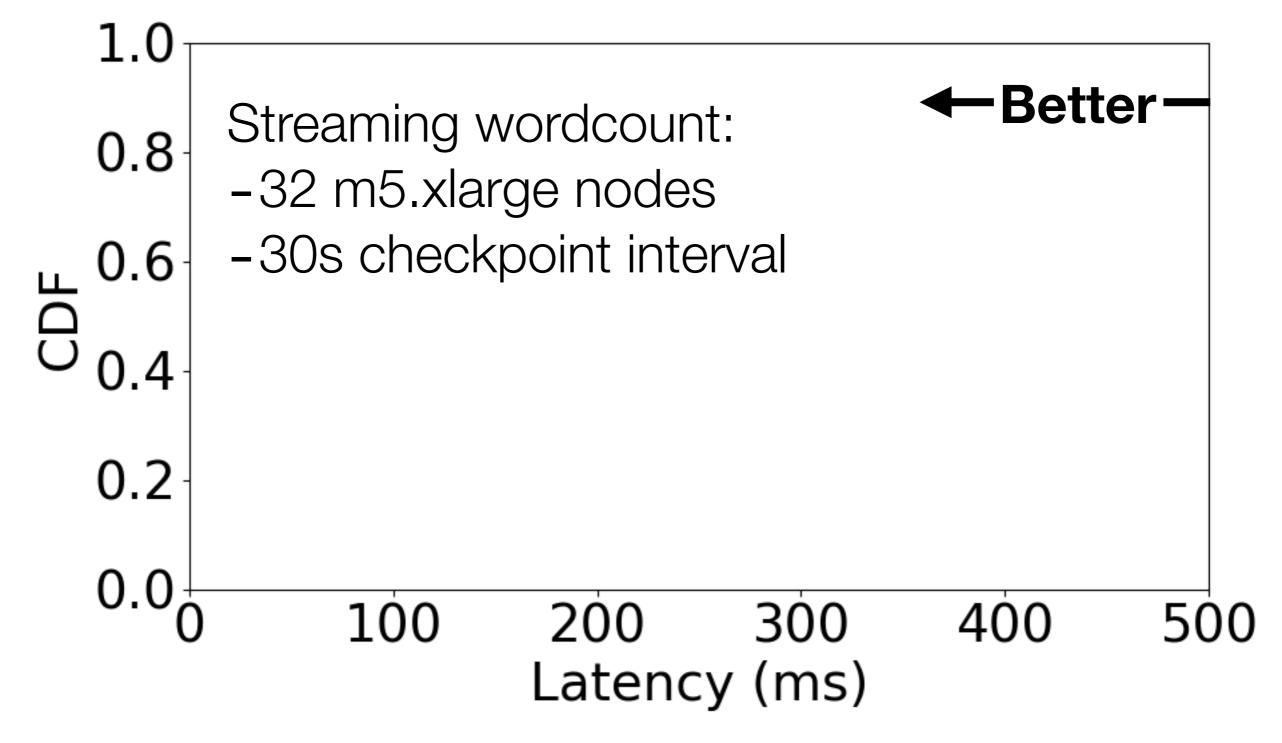


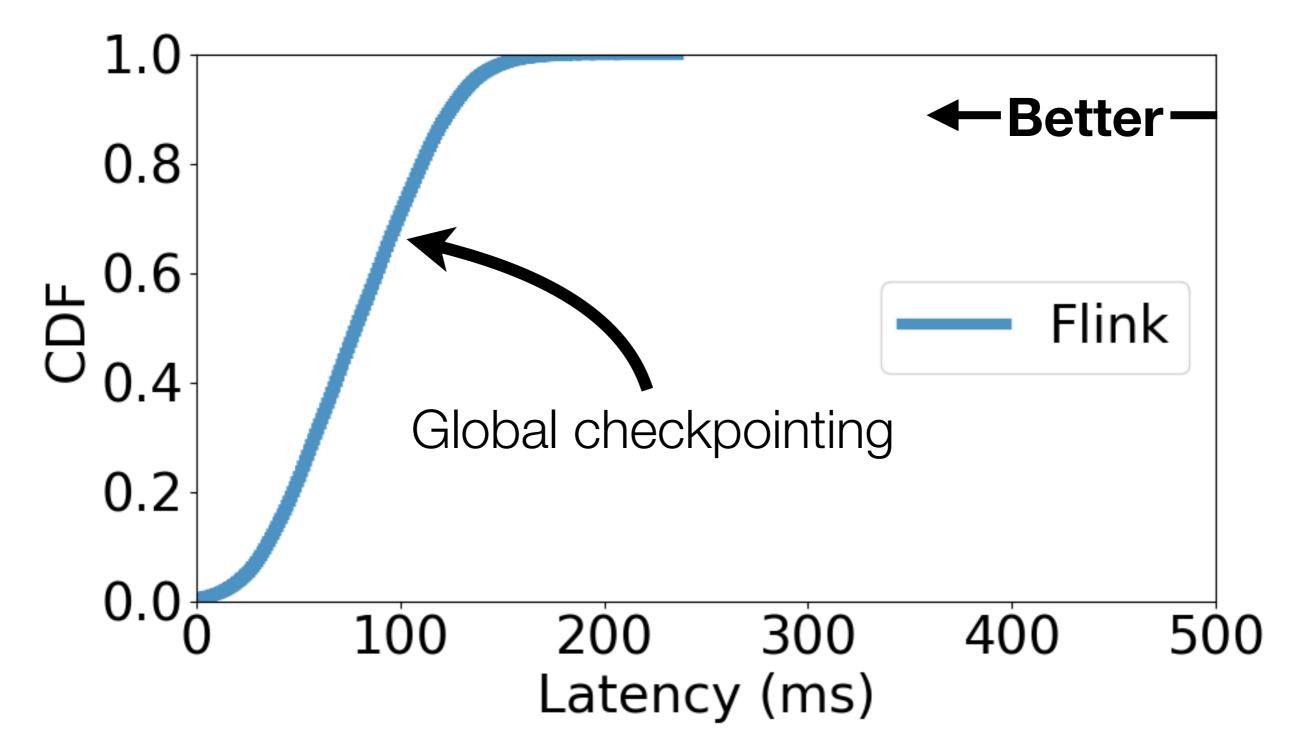


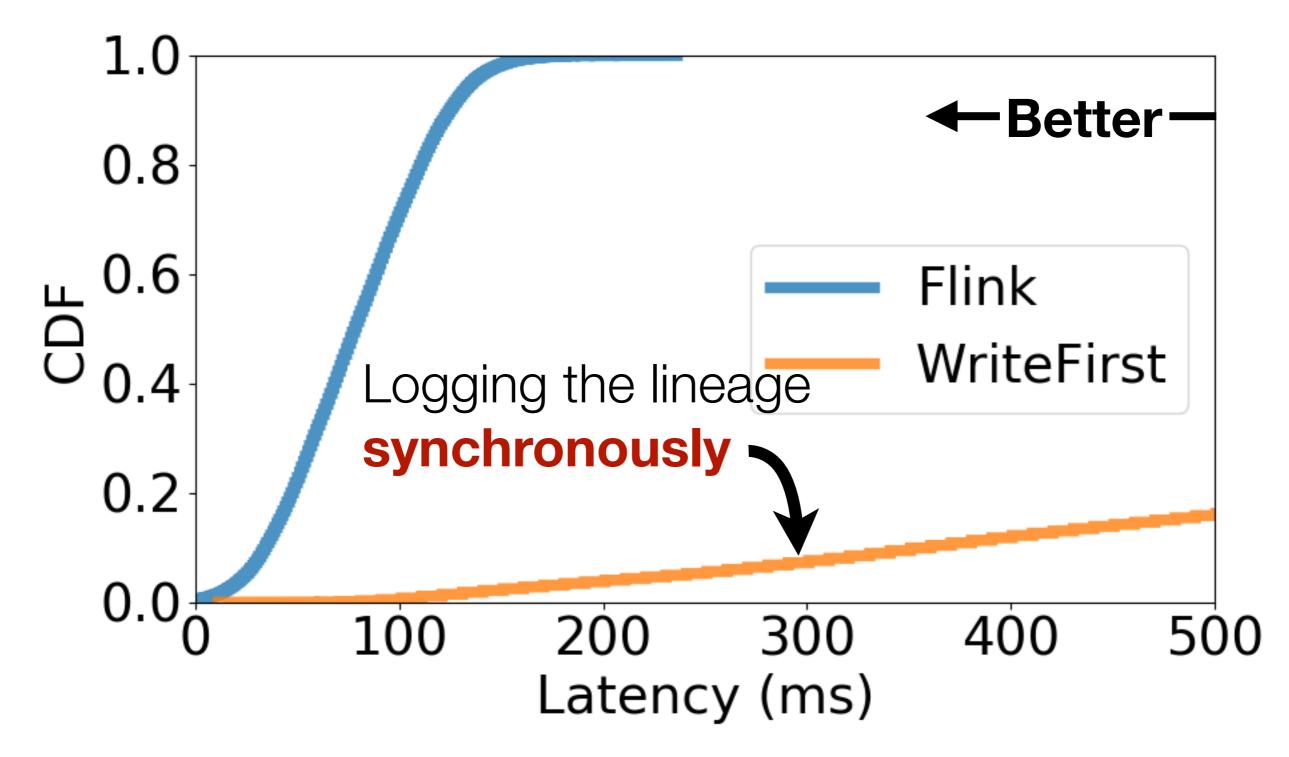
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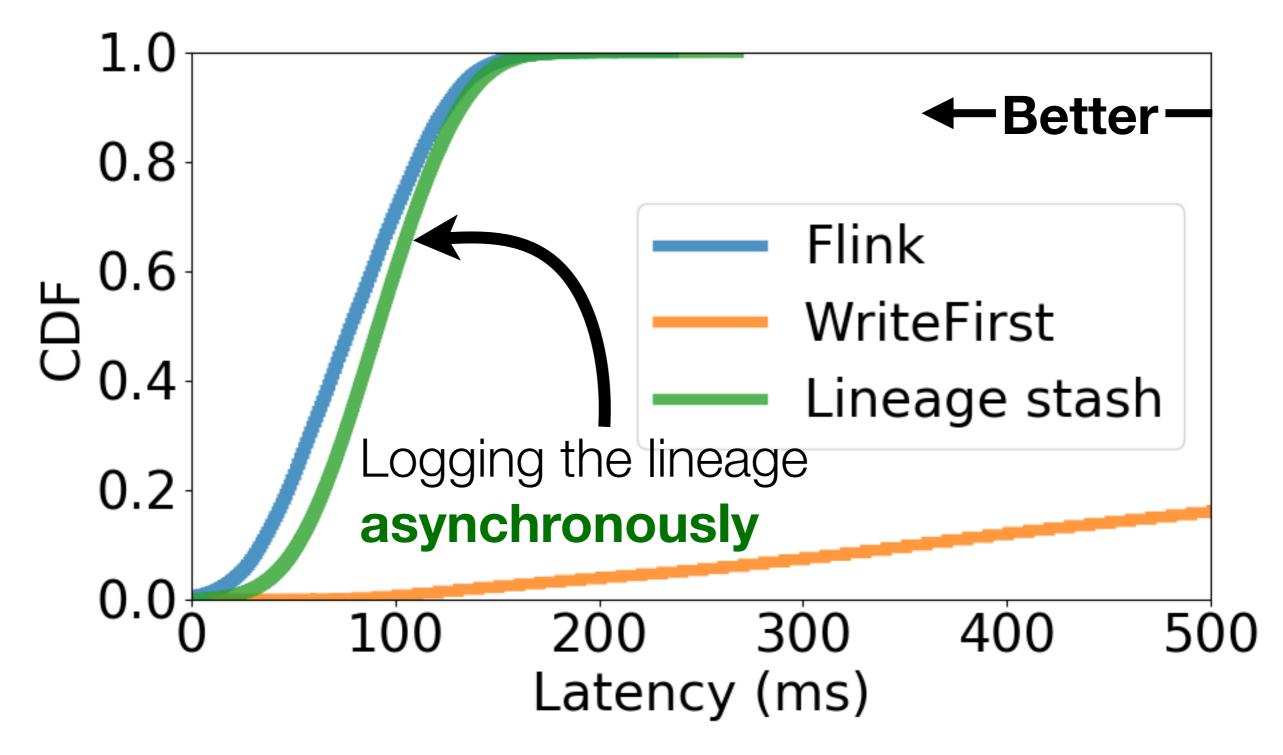


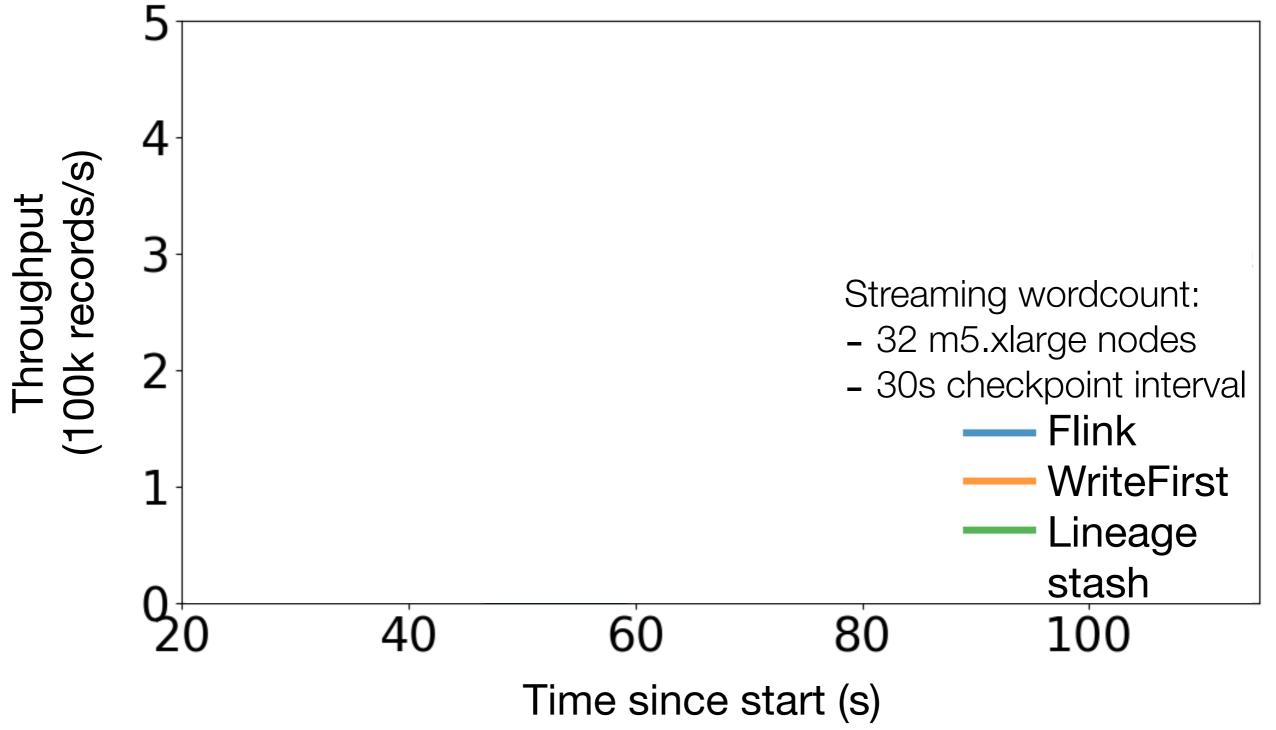
Runtime overhead

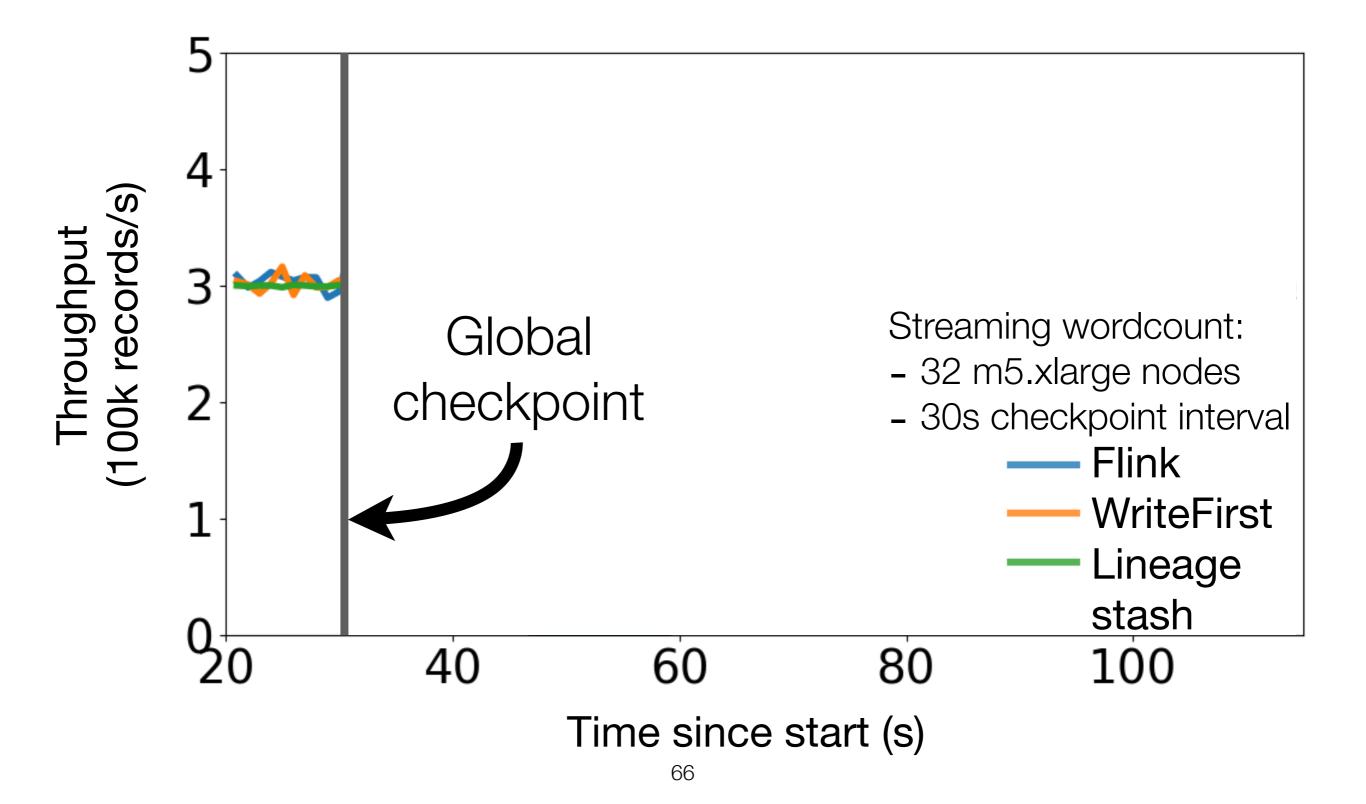


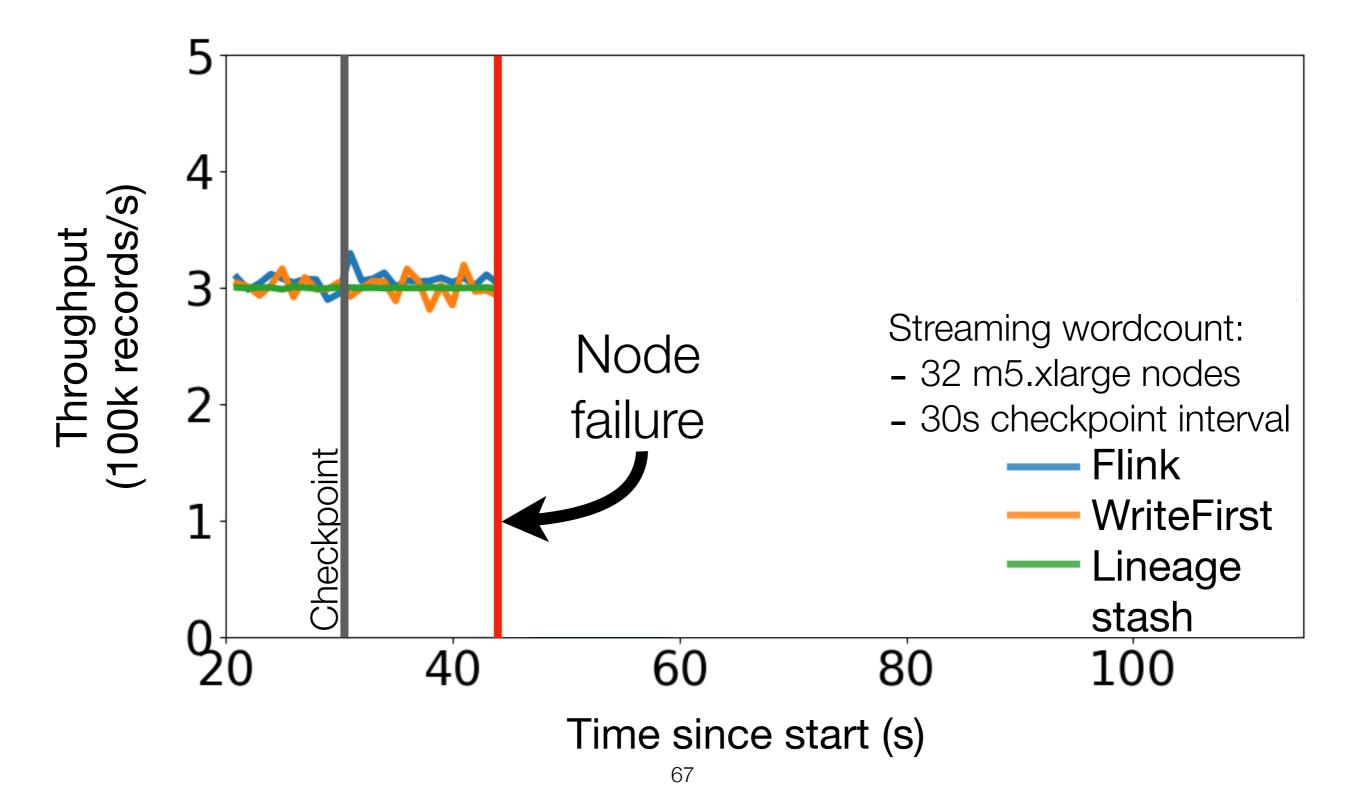


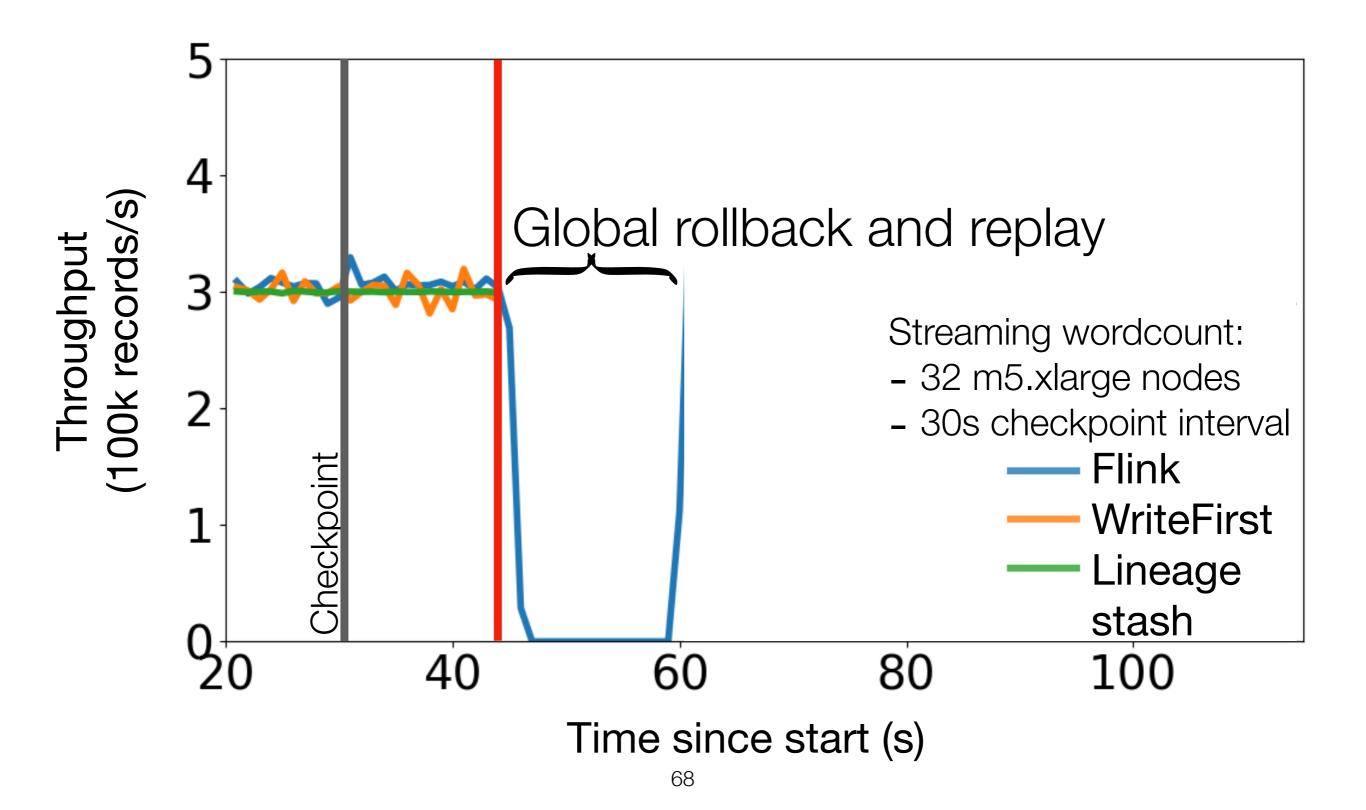


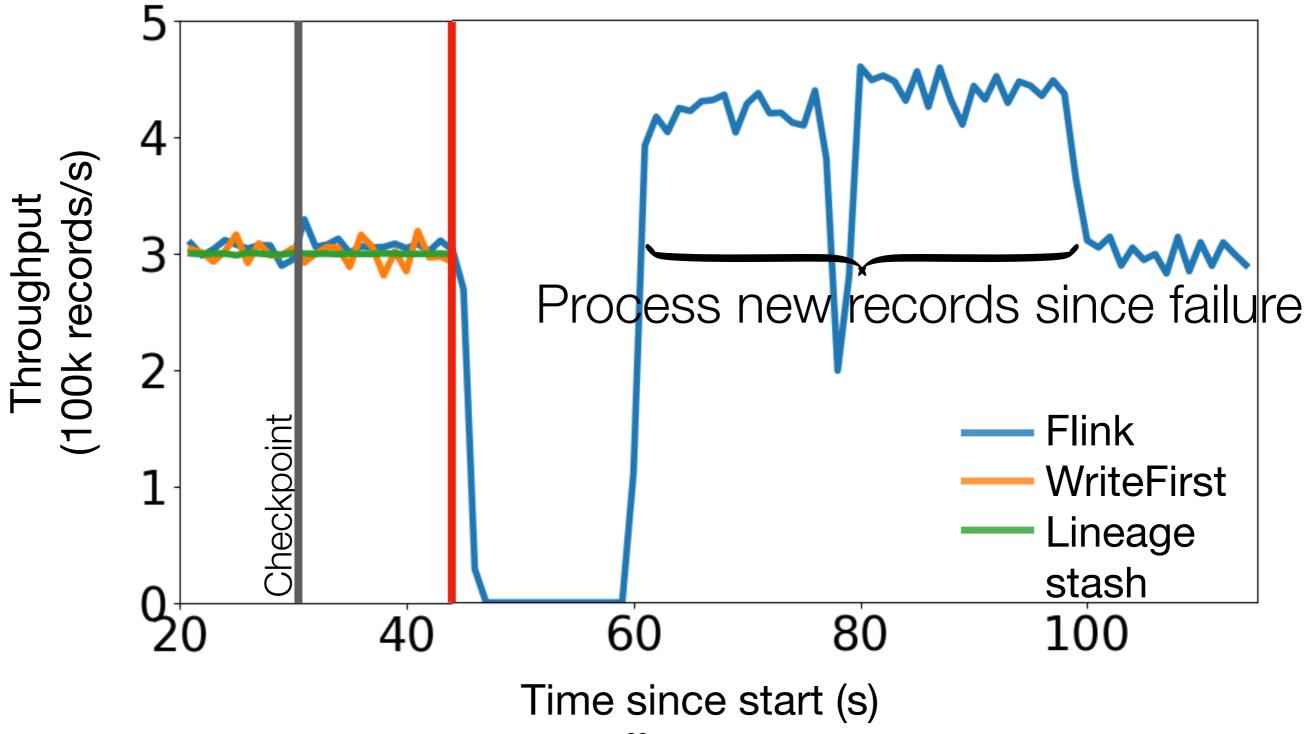


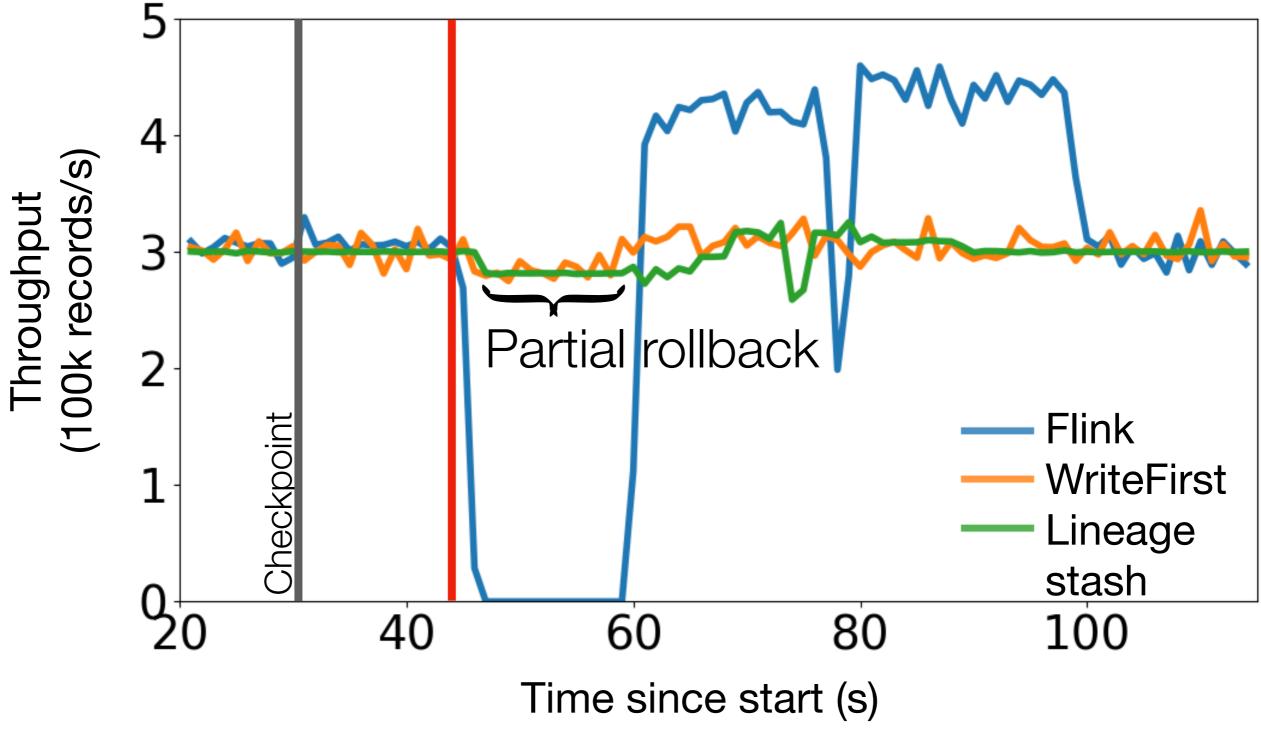




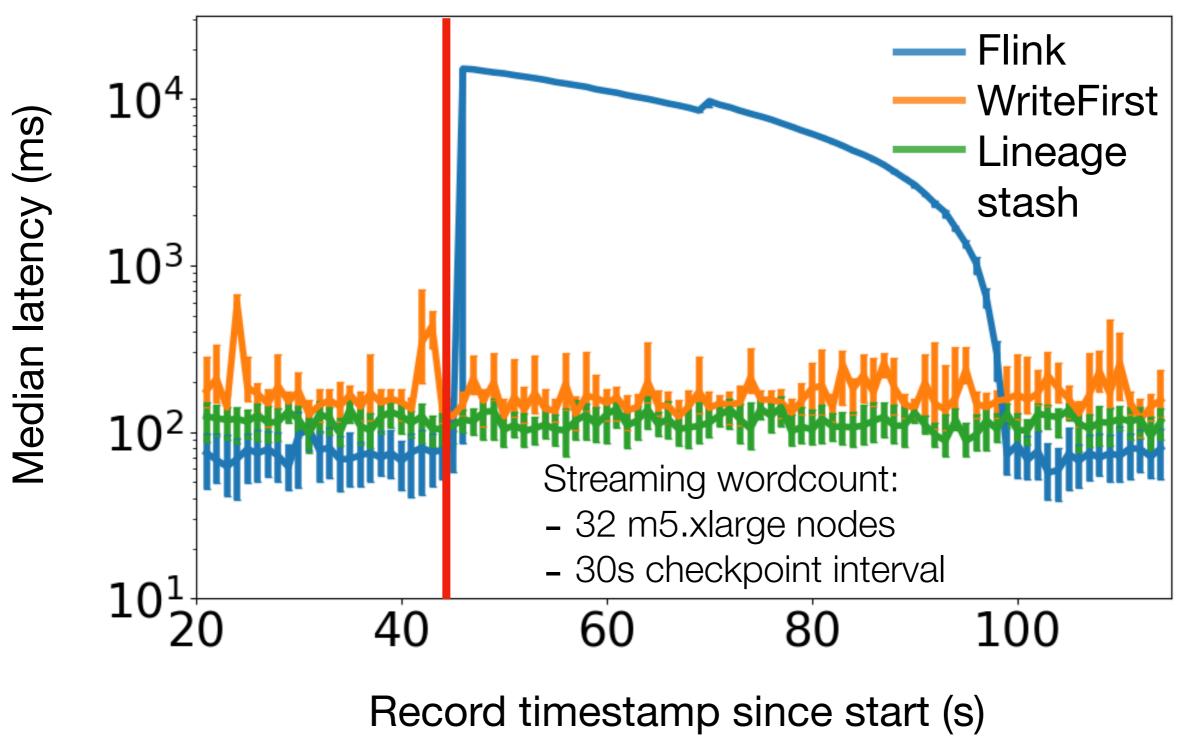








#### Latency during recovery



### Lineage Stash

See the paper (or email me: <a href="mailto:swang@berkeley.edu">swang@berkeley.edu</a>) for:

- Discussion and evaluation of other applications
- Nondeterminism in data processing applications
- Protocols for flushing and recovering the stash

Key idea: **Asynchronously** log the **lineage** and forward **uncommitted** lineage to guarantee recovery correctness.

Low latency during execution and low downtime after a failure for large-scale decentralized data processing applications.